Quantum Sensing and Information-Theoretic Foundations

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Quantum Sensing: Principles and Advantages

$$|\mathbf{1}\rangle \frac{|\mathbf{1}\rangle}{|\mathbf{E}|} = \hbar\omega_0 \qquad |\mathbf{1}\rangle \frac{|\mathbf{1}\rangle}{|\mathbf{0}\rangle} \frac{|\mathbf{1}\rangle}{|\mathbf{1}\rangle} \frac{|\mathbf{$$

- Can detect changes of quantum phenomena (below atomic level)
- Changes often related to electric and magnetic fields (but also gravity, displacement, pressure)
 → imply changes in energy levels, spins, photons
- Can also detect quantum phenomena such as entaglement \rightarrow requires preparation and possibly transmission of entangled states

- Atomic vapor/cold clouds (qubits: spin)
- Solid state spins: NMR sensor, P donor in Si, Semiconductor quantum dots (qubits: nuclear/electron spin)
- Elementary particles: muons, neutrons (qubits: spin)
- Supraconducting circuits (qubits: currents, charges of eigenstates)
- Interferometers (qubits: photons)

Quantum Sensing Applications

- Nuclear magnetic resonance spectrography/imaging
- Quantum magnetic encephalography
- High resolution spectroscopy
- Material testing
- Gravimetry
- Radar
- Thermometry
- Ghost imaging

Ex 1: The Ramsey Sensing Protocol



Ramsey-protocol:

•
$$|\psi_0
angle=rac{1}{\sqrt{2}}(|0
angle+|1
angle)$$

- Wish to estimate phase shift: $|\psi(t)\rangle = \frac{1}{\sqrt{2}}(|0\rangle + e^{-i\omega_0 t} |1\rangle))$
- Prepare for measurement: $|\alpha\rangle = \frac{1}{2} (1 + e^{-i\omega_0 t}) |0\rangle + \frac{1}{2} (1 - e^{-i\omega_0 t}) |1\rangle$
- Readout probability: $p = \sin^2(\omega_0 t/2) = \frac{1}{2}(1 - \cos(\omega_0 t))$

Ex 2: Quantum Illumination for Radar Applications



- Detector: perform measurements to detect entanglement
- Repeat experiment many times and apply classical hypothesis test on the measurement data to decide on presence of obstacle
- Practical challenges: losses, decoherence/storage of idlers

Information-Theoretic Foundations

Basic Notions and Measurements

- The state of a system is described by a density matrix $\rho \succeq 0$ with trace 1
- Operations we can perform on a system (Kraus operators)

$$\rho_{\text{out}} = \sum_{i=1}^{d_{\text{in}} \cdot d_{\text{out}}} \mathsf{K}_i \rho_{\text{in}} \mathsf{K}_i^*, \qquad \text{for } \sum_{i=1}^{d_{\text{in}} \cdot d_{\text{out}}} \mathsf{K}_i^* \mathsf{K}_i = \mathsf{I}$$

• Measurements \rightarrow output is a real value $x \in \mathcal{X}$:

$$\Pr(x) = \operatorname{tr}\left(\sum_{i=1}^{d_{\operatorname{in}}} \mathsf{M}_{x,i}\rho_{\operatorname{in}}\mathsf{M}_{x,i}^*\right) = \operatorname{tr}\left(\sum_{i=1}^{d_{\operatorname{in}}} \mathsf{M}_{x,i}^*\mathsf{M}_{x,i} \cdot \rho_{\operatorname{in}}\right) = \operatorname{tr}(\Lambda_x\rho_{\operatorname{in}})$$

where $\sum_{x} \Lambda_x = I$.

The Symmetric Detection Problem

- If $\mathsf{H}=\mathsf{0}$ then the system has density matrix ρ
- If H = 1 then the system has density matrix σ
- \bullet Apply measurement $\{\Lambda_0,\Lambda_1=I-\Lambda_0\}$ to determine $\hat{H}\in\{0,1\}$
- Type-I and II error prob.: $\alpha = 1 \operatorname{tr}(\Lambda_0 \rho)$ and $\beta = \operatorname{tr}(\Lambda_0 \sigma)$
- Wish to minimize $\alpha + \beta = 1 \operatorname{tr} (\Lambda_0(\rho \sigma))$

Optimal Λ_0^* : Project on positive eigenspace of $\rho - \sigma$ If $\rho - \sigma = \sum_i \eta_i |v_i\rangle\langle v_i|$ then $\Lambda_0^* = \sum_{i: \eta_i > 0} |v_i\rangle\langle v_i|$.and $\alpha^* + \beta^* = 1 - \sum_{i: \eta_i > 0} \eta_i$

The Asymmetric Detection Problem

- $\bullet~{\rm If}~{\sf H}=0~/~{\sf H}=1$ the system has density matrix $\rho~/~\sigma$
- Maximize type-II error prob. under type-I error constraint $\alpha \leq \epsilon :$

$$\beta^{*}(\epsilon) = \inf_{\substack{\Lambda_{0} \succ 0:\\ I - \Lambda_{0} \succeq 0\\ \operatorname{tr}((I - \Lambda_{0})\rho) \leq \epsilon}} \operatorname{tr}(\Lambda_{0}\sigma)$$

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• Semi-definite programme \rightarrow dual problem

$$\beta^{*}(\epsilon) = \sup_{\substack{\mathsf{A} \succeq 0, \mu \ge 0\\ \mathsf{A} \succeq \mu \rho - \sigma}} (\operatorname{tr}(\mu \rho) - \operatorname{tr}(\mathsf{A}) - \mu \epsilon)$$

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• Given μ , optimal A^{*}:

$$\mu \rho - \sigma = \sum_{i} \eta_{i} |\mathbf{v}_{i}\rangle\langle\mathbf{v}_{i}| \qquad \Rightarrow \qquad \mathsf{A}^{*} = \sum_{i: \eta_{i} \ge 0} \eta_{i} |\mathbf{v}_{i}\rangle\langle\mathbf{v}_{i}|$$

The Quantum Neyman-Pearson Test

• Dual Problem:
$$\beta^*(\epsilon) = \sup_{\mu \ge 0} (\operatorname{tr}(\mu\rho) - \operatorname{tr}(\mathsf{A}^*) - \mu\epsilon)$$

for $\mathsf{A}^* = \sum_{i: \ \eta_i \ge 0} \eta_i |\mathbf{v}_i\rangle\langle\mathbf{v}_i|$ for (η_i, \mathbf{v}_i) EVs of $(\mu\rho - \sigma)$

• Complementary Slackness Conditions: $\mu^* \operatorname{tr} ((I - \Lambda_0^*)\rho) = \mu^* \epsilon$ and

$$A^*(I - \Lambda_0^*) = 0$$
 and $(\mu^* \rho - \sigma - A^*)\Lambda_0^* = 0$

Solution:

 Λ_0^* projects onto non-negative eigenspace of $(\mu^* \rho - \sigma)$ for a $\mu^* \ge 0$ that is adjusted in function of type-I constraint ϵ : tr $((I - \Lambda_0^*)\rho) = \epsilon$

The Asymptotic Regime: Stein Exponent

- The asymptotic case with many observations
- $\bullet\,$ If $\mathsf{H}=\mathsf{0}$ / $\mathsf{H}=\mathsf{1}$ the system has density matrix $\rho^{\otimes n}$ / $\sigma^{\otimes n}$

Stein exponent (with Umegaki's relative entropy

$$\lim_{n\to\infty} -\frac{1}{n}\log\beta_n^*(\epsilon) = D(\rho\|\sigma) := \operatorname{tr}\left(\rho(\log\rho - \log\sigma)\right)$$

^[1] F. Hiai and D. Petz, "The Proper Formula for Relative Entropy and its Asymptotics in Quantum Probability," Commun. Math. Phys., 1991.

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• Consider spectral decomposition

$$\rho^{\otimes n} - e^{nt} \sigma^{\otimes n} = \sum_{i} \eta_i \underbrace{|\mathbf{v}_i\rangle\langle\mathbf{v}_i|}_{E_i}$$

• Optimal test is projection on nonnegative eigenspaces $\{E_i\}_{i: \eta_i \ge 0}$

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• Relate to classical test between: $\{p_i = \operatorname{tr} (\rho^{\otimes n} E_i)\}$ and $\{q_i = \operatorname{tr} (\rho^{\otimes n} E_i)\}$

• Define
$$K_n = \sum_{i: \eta_i > 0} E_i$$

• Notice: tr
$$\left(\rho^{\otimes n}K_{n}\right) = \sum_{i: \eta_{i}>0} p_{i} \leq e^{-tsn}\sum_{i} p_{i}^{s+1}q_{i}^{-s}$$

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	v	

$$\begin{aligned} \forall t > 0: \qquad \underbrace{\operatorname{tr}\left(\rho^{\otimes n}\Lambda_{0,n}\right)}_{\geq 1-\epsilon} &= \operatorname{tr}\left(\left(\rho^{\otimes n} - e^{nt}\sigma^{\otimes n}\right)\Lambda_{0,n}\right) + \operatorname{tr}\left(e^{nt}\sigma^{\otimes n}\Lambda_{0,n}\right) \\ &\leq \operatorname{tr}\left(\left(\rho^{\otimes n} - e^{nt}\sigma^{\otimes n}\right)K_{n}\right) + \operatorname{tr}\left(e^{nt}\sigma^{\otimes n}\Lambda_{0,n}\right) \\ &\leq \operatorname{tr}\left(\rho^{\otimes n}K_{n}\right) + e^{nt}\underbrace{\operatorname{tr}\left(\sigma^{\otimes n}\Lambda_{0,n}\right)}_{\beta_{n}} \end{aligned}$$

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• Thus, finally: $-\frac{1}{n}\log\beta_n \le t - \frac{1}{n}\log(1 - \epsilon - e^{-n\varphi(t)})$ for $\varphi(t) := \sup_{0\le s\le 1} \left(ts - \log \operatorname{tr}\left(\rho^{1+s}\sigma^{-s}\right)\right)$ and $\phi(t) > 0$ for $t \ge D(\rho \| \sigma)_{13/25}$

The Asymptotic Regime: Chernoff-Hoeffding Exponents

• If $\mathsf{H}=\mathsf{0}$ / $\mathsf{H}=\mathsf{1}$ the system has density matrix $\rho^{\otimes n}$ / $\sigma^{\otimes n}$

Chernoff exponent: Maximum symmetric exponent

$$\sup_{\{\Lambda_{0,n}\}} \lim_{n \to \infty} -\frac{1}{n} \log(\alpha_n + \beta_n) = \sup_{0 \le s \le 1} -\log \operatorname{tr} \left(\sigma^s \rho^{1-s}\right)$$

^[3] H. Nagaoka, "The converse part of the theorem for quantum Hoeffding bound," Arxiv, Nov. 2006
[4] K.M.R. Audenaert et al., "The Quantum Chernoff Bound," Phys. Review Letters 2007
[5] M. Hayashi, "Error exponent in asymmetric quantum hypothesis testing and its application to classical-quantum channel coding," Physical Review A, 2007.

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• Asymmetric Hoeffding exp: $\beta_{CH,n}(r) = \inf_{\substack{\Lambda_0:\\ \operatorname{tr}((I-\Lambda_0)\rho^{\otimes n}) \leq e^{-rn}}} \operatorname{tr}(\Lambda_0 \sigma^{\otimes n})$



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Proof Sketch of Converse to the Chernoff Bound

- Spectral decompositions $\rho = \sum_{i} \eta_i |x_i\rangle \langle x_i|$ and $\sigma = \sum_{j} \nu_j |y_j\rangle \langle y_j|$
- For any projectors $\Lambda,$ (i.e., $\Lambda\Lambda=\Lambda):$

$$\beta = \operatorname{tr} (\Lambda \sigma) = \operatorname{tr} (\Lambda | \Lambda \sigma) = \sum_{i} \operatorname{tr} (\Lambda | x_i \rangle \langle x_i | \Lambda \sigma)$$
$$= \sum_{j,i} \nu_j \operatorname{tr} (\Lambda | x_i \rangle \langle x_i | \Lambda | y_j \rangle \langle y_j |) = \sum_{j,i} \nu_j |\langle x_i | \Lambda | y_j \rangle|^2$$

• Similarly: $\alpha = \sum_{j,i} \eta_i |\langle x_i | (I - \Lambda) | y_j \rangle|^2$

$$\alpha + \beta = \sum_{j,i} \eta_i |\langle x_i | (I - \Lambda) | y_j \rangle|^2 + \sum_{j,i} \nu_j |\langle x_i | \Lambda | y_j \rangle|^2$$

Proof Sketch of Converse to the Chernoff Bound II

Reduce to classical test using $\eta |u - v|^2 + \nu |v|^2 \ge \frac{1}{2} |u|^2 \min\{\eta, \nu\}$:

$$\begin{aligned} \alpha + \beta &= \sum_{j,i} \eta_i \left| \langle x_i | \mathbf{I} | y_j \rangle - \langle x_i | \Lambda | y \rangle_j \right|^2 + \sum_{j,i} \nu_j \left| \langle x_i | \Lambda | y_j \rangle \right|^2 \\ &\geq \frac{1}{2} \sum_{j,i} \min\{\eta_i, \nu_j\} |\langle x_i | y_j \rangle|^2 = \frac{1}{2} \sum_{j,i} \min\{p_{ij}, q_{ij}\} \end{aligned}$$

• Cannot do better than optimal classical test between $\{p_{ij} = \eta_i \langle x_i | y_j \rangle\}$ and $\{q_{ij} = \nu_i \langle x_i | y_j \rangle\}$:

$$\Rightarrow \alpha_n + \beta_n \geq \frac{1}{2} P^*_{\text{class-HT}}(\boldsymbol{p}^n, \boldsymbol{q}^n)$$

• Use equality
$$-\log \operatorname{tr} \left(\sigma^{s} \rho^{1-s} \right) = -\log \sum_{i,j} q_{ij}^{s} p_{ij}^{1-s}$$

Distributed Detection with Positive Rates (Rouzé et al.)



• Quantum testing against independence: $\otimes n$

$$\begin{split} \mathsf{H} &= 0 \colon \qquad \rho_{AB}^{\otimes n} \\ \mathsf{H} &= 1 \colon \qquad \tilde{\rho}_{AB}^{\otimes n} = \rho_{A}^{\otimes n} \otimes \rho_{B}^{\otimes n} \end{split}$$

- Communication rate constraint $\frac{1}{n} \log |\mathcal{W}_n| \le R$; $(B_n \text{ local obs.})$
- Largest Stein exponent: $\sup_{\{T_n\}} \lim_{n \to \infty} -\frac{1}{n} \log \Pr[\hat{H} = 0 | H = 1]$

s.t. $\lim_{n\to\infty} \alpha_n = \Pr[\hat{H} = 1 | H = 0] = 0$ (vanishing)

Distributed Detection with Positive Rates (Rouzé et al.)



• Classical-quantum setup:

$$\begin{aligned} \mathsf{H} &= 0: \qquad \left(\sum_{a} P(a) |a\rangle \langle a| \otimes \rho_{B,a} \right)^{\otimes n} \\ \mathsf{H} &= 1: \qquad \left(\sum_{a} \tilde{P}(a) |a\rangle \langle a| \right)^{\otimes n} \otimes \rho_{B}^{\otimes n} \end{aligned}$$

• Largest Stein exponent (for all ϵ):

$$\sup_{\{T_n\}} \lim_{n \to \infty} -\frac{1}{n} \log \Pr[\hat{H} = 0 | H = 1]$$

s.t. $\lim_{n \to \infty} \alpha_n = \Pr[\hat{H} = 1 | H = 0] = \epsilon \in [0, 1)$

Distributed Detection with Zero Rates (Sreekumar et al.)



- Stein exponent at the receiver when B^n is observed locally
- Solved in terms of single-letter expression only in very special cases
- When only classical communication is allowed, a single bit suffices
- Classical communication can be infinitely worse than sending a single qubit

Distributed Detection: Some Open Problems

- Optimal exponents in fully quantum case?
- Multiple sensors or network scenarios?
- Noisy channels? (Dichotomy in the classical setup)
- Variable-rate coding?

Integrated Sensing and Communications (ISAC)



- \bullet Sensing: Detection exponent for θ
- Input signal influences both decoding and sensing performance



• Recent works: other sensing performances (estimation) and benefits of entanglement

[6] S.-Y. Wang et al., "Joint Quantum Communication and Sensing," ITW 2022.

A Few Words on Parameter Estimation

• There exists a Quantum Cramer-Rao Bound (QCRB) :

$$\begin{split} \operatorname{MSE}(\hat{\theta}) &\geq \frac{1}{\operatorname{tr}\left(\mathcal{R}_{\rho\theta}^{-1}\left(\frac{\delta\rho_{\theta}}{\delta\theta}\right)\rho_{\theta}\mathcal{R}_{\rho\theta}^{-1}\left(\frac{\delta\rho_{\theta}}{\delta\theta}\right)\right)} \\ \text{for } \mathcal{R}_{\rho}^{-1}(\mathcal{O}) &= \sum_{j,k} \frac{2}{\eta_{j} + \eta_{k}} \langle j | \mathcal{O} | k \rangle | j \rangle \langle k | \text{ and } \rho = \sum_{j} \eta_{i} | i \rangle \langle i | \rangle \end{split}$$

- Challenge is then to prepare a good state ρ and provide measurements "exhausting" the QCRB
- Requirement for preparing entangled states
- In certain scenarios, measurements can depend on parameter \rightarrow Precision beyond the QCRB...

- Quantum sensing/quantum metrology is an exciting field with decades of great success and huge interest at the moment
- New applications seem more imminent than for other quantum fields
- Theoretical foundations still bear lots of fundamental open questions and challenges
- Quantum nature with states in different eigenbases make information theoretic problems hard
 - \rightarrow still constant progress on the topic

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