

Side-Information with a Grain of Salt

Pliny's Naturalis Historia, 77 A.D.:

Take two dried walnuts, two figs, and twenty leaves of rue; pound them all together, with the addition of a grain of salt; if a person takes this mixture fasting, he will be proof against all poisons for that day.

Michèle Wigger

ETH Zurich
wigger@isi.ee.ethz.ch

Telecom ParisTech, 1 April 2009

ETH

Eidgenössische Technische Hochschule Zürich
Swiss Federal Institute of Technology Zurich

Signal and Information
Processing Laboratory
Institut für Signal- und
Informationsverarbeitung



*Based on collaborations with
Shraga Bross, Michael Gastpar, Gerhard Kramer, Amos Lapidoth, Shlomo Shamai*

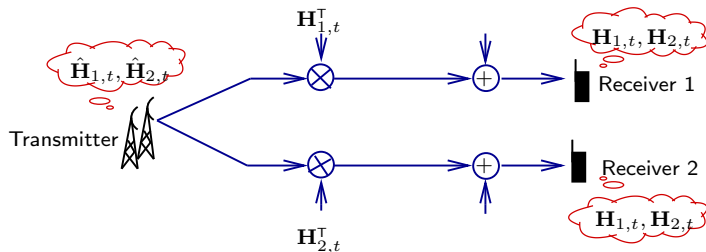
Robustness w.r.t. Precision of Side-Information

- ▶ Performance with imprecise side-information
- ▶ Robust schemes

Four Communication Scenarios with Side-Information

A on channel states in MIMO broadcast channels

Fading MIMO BC with Imprecise State-Information



Lapidoth/Shamai/Wigger'05

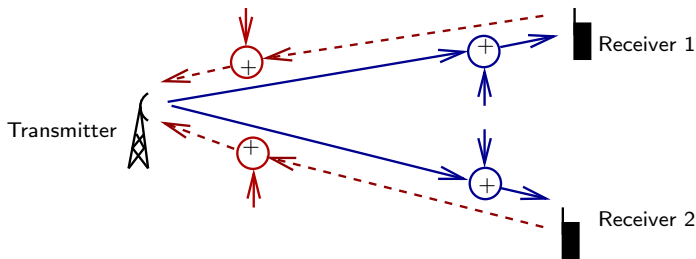
Gain (i.e., 2 efficient links thanks to two antennas)
collapses with imprecise transmitter state-information!

Four Communication Scenarios with Side-Information

A on channel states in MIMO broadcast channels

B via feedback in Gaussian broadcast channels

Scalar Gaussian Broadcast Channel with Noisy Feedback



Gastpar/Lapidoth/Wigger'09

Gain collapses completely when feedback is noisy!
I.e. # efficient links collapses from 2 to 1!

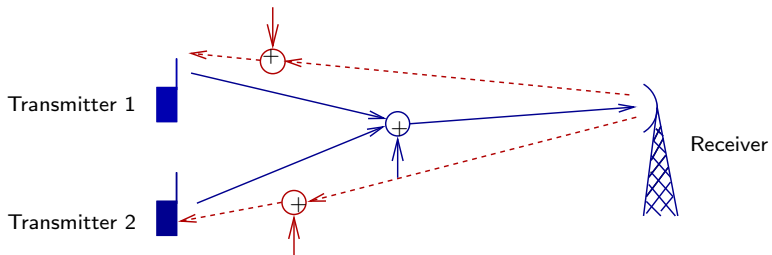
Four Communication Scenarios with Side-Information

A on channel states in MIMO broadcast channels

B via feedback in Gaussian broadcast channels

C via feedback in Gaussian multiple-access channels

Gaussian MAC with Noisy Feedback



Lapidoth&Wigger'06

- ▶ Noisy feedback is almost as good as perfect feedback!
- ▶ Even if noisy, feedback is *always* beneficial!

Four Communication Scenarios with Side-Information

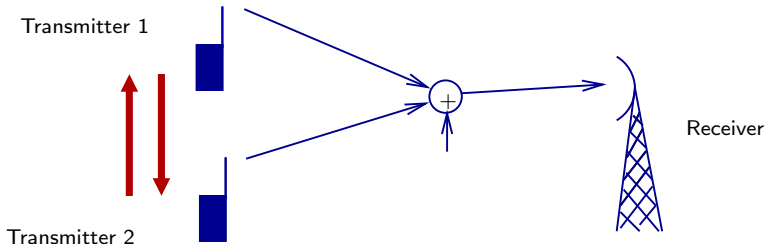
A on channel states in MIMO broadcast channels

B via feedback in Gaussian broadcast channels

C via feedback in Gaussian multiple-access channels

D on other transmitter's message in Gaussian multiple-access channels

Gaussian MAC with Conferencing Encoders à la Willems'83



Bross/Lapidoth/Wigger'08

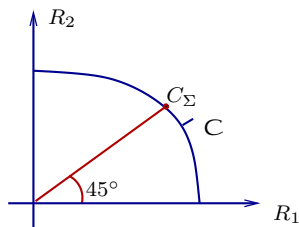
- ▶ Capacity region
- ▶ Best conferencing: optimally describe message to other transmitter

Four Communication Scenarios with Side-Information

- A on channel states in MIMO broadcast channels
- B via feedback in Gaussian broadcast channels
- C via feedback in Gaussian multiple-access channels
- D on other transmitter's message in Gaussian multiple-access channels

Capacity Region

- ▶ Rates of communication R_1 and R_2
- ▶ **Capacity region C** : Set of (R_1, R_2) s.t. $p(\text{error})$ arbitrarily small
- ▶ **Sum-rate capacity C_Σ** : maximum throughput s.t. $p(\text{error})$ arbitrarily small



Degrees of Freedom η

$$C_{\Sigma}(P) \approx \eta \cdot \frac{1}{2} \log \left(1 + \frac{P}{N} \right), \quad \frac{P}{N} \gg 1.$$

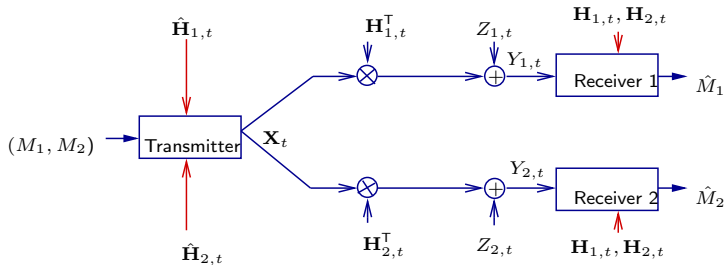
Engineering intuition:

- ▶ # interference-free Gaussian channels in a system
- ▶ $\eta \leq \min\{\# \text{ tx-antennas}, \# \text{ rx-antennas}\}$

Part A

Fading MIMO Gaussian Broadcast Channel with Imprecise State-Information

Gaussian Fading MIMO Broadcast Channel

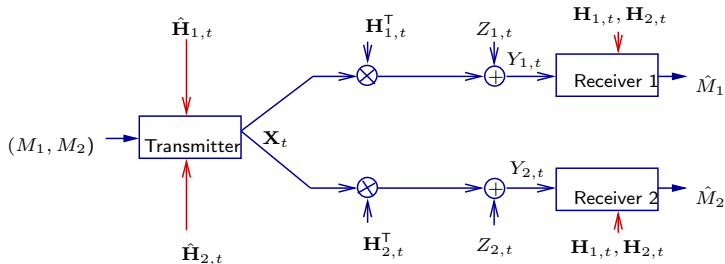


▶ Transmitter has 2 antennas, receivers 1 antenna

▶ $Y_{\nu,t} = \mathbf{H}_{\nu,t}^T \mathbf{X}_t + Z_{\nu,t}; \quad \{Z_{\nu,t}\} \text{ IID } \sim \mathcal{N}(0, N)$

▶ Power constraint: $\frac{1}{n} \sum_{t=1}^n \mathbf{E} [\|\mathbf{X}_t\|^2] \leq P$

Gaussian Fading MIMO Broadcast Channel



- ▶ $\mathbf{H}_{\nu,t} = \hat{\mathbf{H}}_{\nu,t} + \tilde{\mathbf{H}}_{\nu,t}$
- ▶ Tx knows a-causally *realizations* of $\{\hat{\mathbf{H}}_{1,t}\}, \{\hat{\mathbf{H}}_{2,t}\}$
- ▶ Rxs know *realizations* of $\{\tilde{\mathbf{H}}_{1,t}\}, \{\tilde{\mathbf{H}}_{2,t}\}, \{\hat{\mathbf{H}}_{1,t}\}, \{\hat{\mathbf{H}}_{2,t}\}$ (optimistic)

Perfect Channel State Information (CSI)

- ▶ 2 transmit antennas [Caire & Shamai'03]: $\eta_{\text{PerfectCSI}} = 2$
- ▶ 1 transmit antenna: $\eta_{\text{PerfectCSI}} = 1$
- ▶ \implies 2 antennas double throughput at high powers!
- ▶ Beamforming: Transmission in two orthogonal directions

Approximate Channel State Information

Theorem 1: (Lapidoth/Shamai/Wigger'05)

Degrees of freedom collapse from 2 to at most

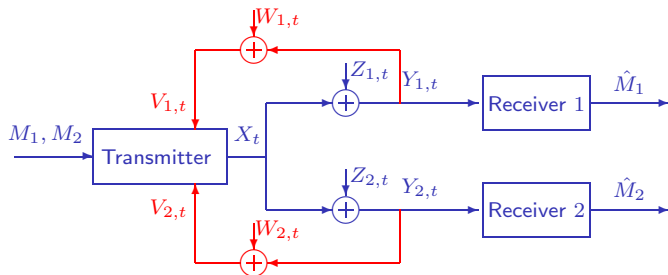
$$\eta_{\text{Approx. CSI}} \leq \frac{4}{3}$$

- ▶ Conjecture: 1 degree of freedom! \Rightarrow “No gain” from 2 tx-antennas!
- ▶ Here: precision of transmitter CSI fixed, not improved as $P \rightarrow \infty$

Part B

Scalar Gaussian Broadcast Channel with Noisy Feedback

Scalar Gaussian Broadcast Channel with Noisy Feedback

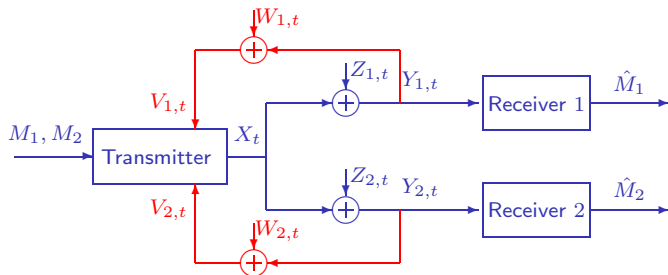


- ▶ Transmitter & receivers: each 1 antenna!

- ▶ $Y_{\nu,t} = X_t + Z_{\nu,t}, \quad \begin{pmatrix} Z_{1,t} \\ Z_{2,t} \end{pmatrix} \sim \mathcal{N}\left(\mathbf{0}, \begin{pmatrix} N & \rho_z N \\ \rho_z N & N \end{pmatrix}\right)$

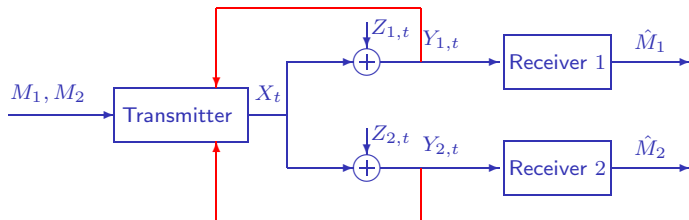
- ▶ Noise correlation caused by interference

Scalar Gaussian Broadcast Channel with Noisy Feedback



- ▶ $V_{\nu,t} = Y_{\nu,t} + W_{\nu,t}$, $\left\{ \begin{pmatrix} W_{1,t} \\ W_{2,t} \end{pmatrix} \right\}$ IID $\sim \mathcal{N}\left(\mathbf{0}, \begin{pmatrix} \sigma_1^2 & 0 \\ 0 & \sigma_2^2 \end{pmatrix}\right)$
- ▶ $\{W_{1,t}, W_{2,t}\}$ independent of $\{Z_{1,t}, Z_{2,t}\}$
- ▶ $X_t = f_t(M_1, M_2 \mathbf{V}_{1,1}^{t-1}, \mathbf{V}_{2,1}^{t-1})$
- ▶ Power constraint: $\frac{1}{n} \sum_{t=1}^n \mathbb{E}[X_t^2] \leq P$

Perfect Feedback



Theorem 2: (Gastpar&Wigger'08)

$$\eta_{\text{PerfectFB}} = \begin{cases} 1, & -1 < \rho_z \leq 1, \\ 2, & \rho_z = -1. \end{cases}$$

For $\rho_z = -1$:

- ▶ 2 degrees of freedom with 1 transmit antenna!
- ▶ Perfect feedback doubles degrees of freedom

Nosiy Feedback: Degrees of Freedom Collapse

Theorem 3: (Gastpar/Lapidoth/Wigger'09, in preparation)

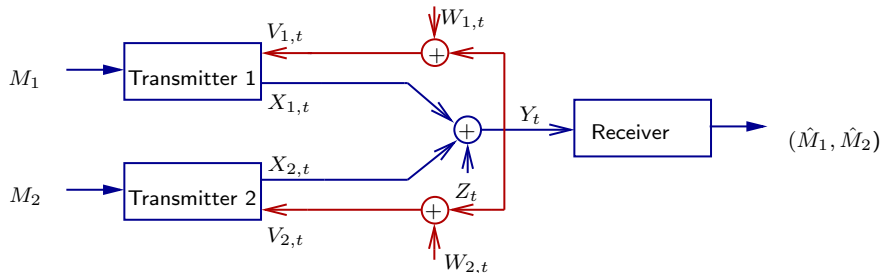
$$\eta_{\text{NoisyFB}} = 1, \quad \text{for all } \rho_z \in [-1, 1].$$

- ▶ Feedback noise variances $\sigma_1^2, \sigma_2^2 > 0$ fixed
- ▶ Gain of [Gastpar/Wigger'08] collapses!
- ▶ Engineering intuition “ $\eta \leq \min\{\#\text{tx-antennas}, \#\text{rx-antennas}\}$ ” OK!

Part C

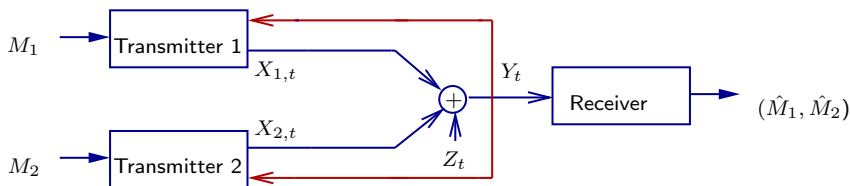
Gaussian MAC with Noisy Feedback

Gaussian MAC with Noisy Feedback



- ▶ $Y_t = X_{1,t} + X_{2,t} + Z_t, \quad \{Z_t\} \text{ IID } \sim \mathcal{N}(0, N)$
- ▶ $V_{\nu,t} = Y_t + W_{\nu,t}, \quad \{(W_{1,t}, W_{2,t})^\top\} \sim \text{IID } \mathcal{N}(\mathbf{0}, \mathbf{K}_{W_1 W_2})$
- ▶ Power constraints: $\frac{1}{n} \sum_{t=1}^n \mathbb{E}[X_{\nu,t}^2] \leq P_\nu, \quad \nu \in \{1, 2\}$

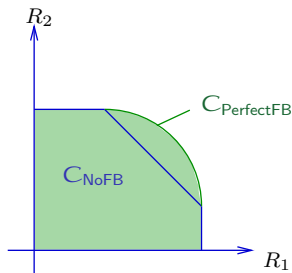
Perfect Feedback



- Ozarow'84:

$$C_{\text{PerfectFB}} =$$

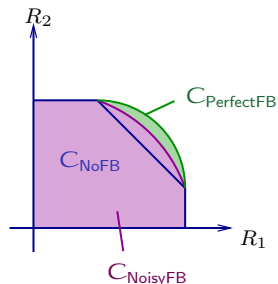
$$\bigcup_{\rho \in [0,1]} \left\{ (R_1, R_2) : \begin{cases} R_1 \leq \frac{1}{2} \log \left(1 + \frac{P_1(1-\rho^2)}{N} \right) \\ R_2 \leq \frac{1}{2} \log \left(1 + \frac{P_2(1-\rho^2)}{N} \right) \\ R_1 + R_2 \leq \frac{1}{2} \log \left(1 + \frac{P_1 + P_2 + 2\sqrt{P_1 P_2} \rho}{N} \right) \end{cases} \right\}$$



- Capacity increased!

Imperfect Feedback: Noisy Feedback

- ▶ C_{NoisyFB} open
- ▶ Ozarow's strategy doesn't work!



- ▶ a) Carleial'82, b) Willems/van der Meulen/Schalkwijk'83, c) Gastpar'05
 - ▶ Collapse to C_{NoFB} if feedback too noisy!
 - ▶ a)&b) do not approach $C_{\text{PerfectFB}}$ when feedback becomes noise-free

Shortcomings of schemes or inherent in problem?

Results for Noisy Feedback

- ▶ We robustify Ozarow's scheme to noisy feedback

Theorem 4: (Lapidoth&Wigger'06)

Noisy-feedback capacity converges to perfect-feedback capacity

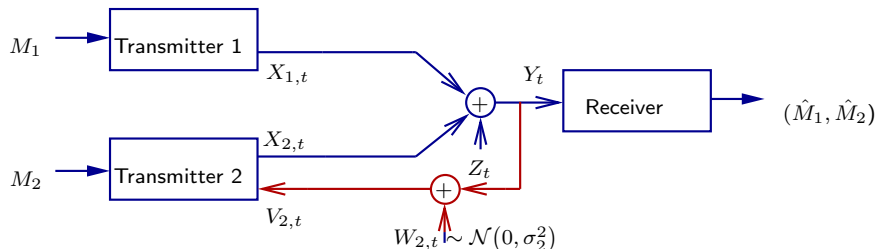
$$\text{cl}\left(\bigcup_{\sigma^2 \geq 0} \bigcap_{\mathbf{K}: \text{tr}(\mathbf{K}) \leq \sigma^2} C_{\text{NoisyFB}}(P_1, P_2, N, \mathbf{K})\right) = C_{\text{PerfectFB}}(P_1, P_2, N)$$

Theorem 5a: (Lapidoth&Wigger'06)

*Noisy feedback is **always** beneficial!*

$$C_{\text{NoFB}}(P_1, P_2, N) \subsetneq C_{\text{NoisyFB}}(P_1, P_2, N, \mathbf{K}_{W_1 W_2}) \quad \forall \mathbf{K}_{W_1 W_2} \succeq 0.$$

Imperfect Feedback: Noisy Partial Feedback



- ▶ Situation even worse! → Feedback beneficial?
- ▶ If $\sigma_2^2 = 0$, Cover-Leung region achievable

Van der Meulen'87:

Is Cover-Leung region capacity when $\sigma_2^2 = 0$?

Results for Partial Feedback

- ▶ Our robust scheme still works!

Theorem 5b: (Lapidoth&Wigger'06)

Noisy partial feedback is *always* beneficial!

$$C_{\text{NoFB}}(P_1, P_2, N) \subsetneq C_{\text{NoisyPartialFB}}(P_1, P_2, N, \sigma_2^2), \quad \forall \sigma_2^2 \geq 0.$$

Theorem 6: Answer to van der Meulen'87 (Lapidoth&Wigger'06)

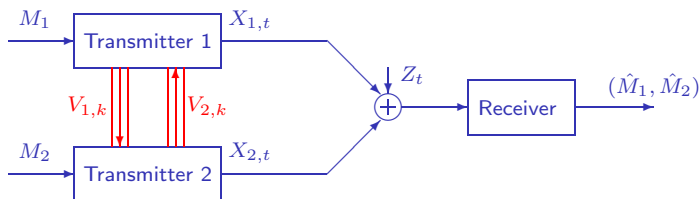
Perfect partial-feedback capacity \neq Cover-Leung region!

$$\mathcal{R}_{\text{CL}}(P_1, P_2, N) \subsetneq C_{\text{PerfectPartialFB}}(P_1, P_2, N), \quad \text{for some } P_1, P_2, N > 0.$$

Part D

Gaussian MAC with Conferencing Encoders

Gaussian MAC with Conferencing Encoders

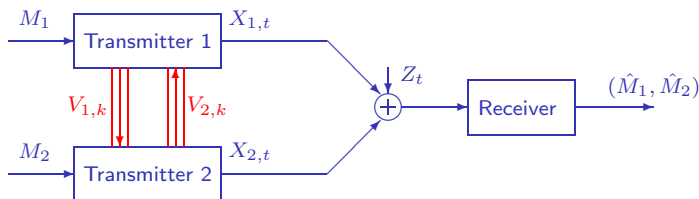


1. phase: Conference (Willems'83)

- ▶ κ sequential uses of perfect bit-pipes
- ▶ $V_{1,k} = \varphi_{1,k}(M_1, V_{2,1}^{k-1})$; $V_{2,k} = \varphi_{2,k}(M_2, V_{1,1}^{k-1})$
- ▶ Rate-limitations:

$$\sum_{k=1}^{\kappa} \log |\mathcal{V}_{1,k}| \leq nC_{12} \quad \text{and} \quad \sum_{k=1}^{\kappa} \log |\mathcal{V}_{2,k}| \leq nC_{21}$$

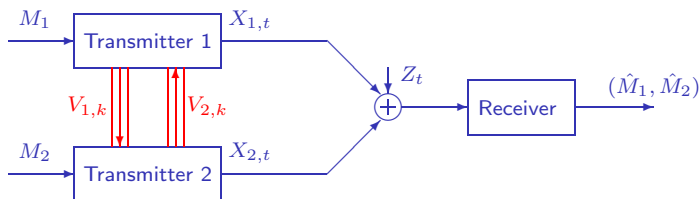
Gaussian MAC with Conferencing Encoders



2. phase: Transmission over channel:

- ▶ $Y_t = X_{1,t} + X_{2,t} + Z_t; \quad \{Z_t\} \text{ IID } \sim \mathcal{N}(0, N)$
- ▶ $X_{1,t} = f_{1,t}(M_1, V_{2,1}^\kappa); \quad X_{2,t} = f_{2,t}(M_2, V_{1,1}^\kappa)$
- ▶ Power constraints: $\frac{1}{n} \sum_{t=1}^n \mathbf{E}[X_{\nu,t}^2] \leq P_\nu, \quad \nu \in \{1, 2\}$

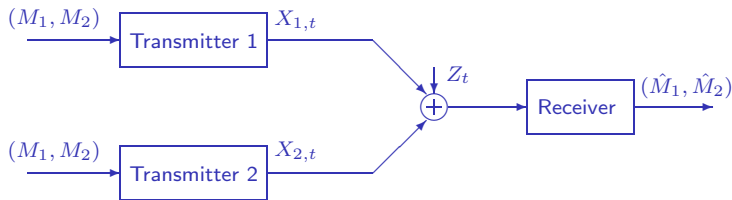
Gaussian MAC with Conferencing Encoders



Special Cases:

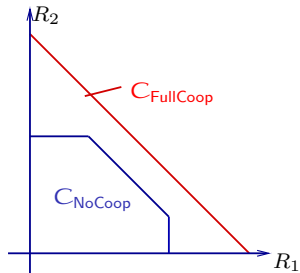
- ▶ $C_{12} = C_{21} = \infty$: full cooperation (both txs know (M_1, M_2))
- ▶ $C_{12} = 0, C_{21} = \infty$: Tx 1 knows (M_1, M_2) , Tx 2 only M_2
- ▶ $C_{12} = C_{21} = 0$: no conferencing

Full Cooperation



- ▶ Transmitters 1 and 2 know (M_1, M_2) :

$$C_{\text{FullCoop}} = \left\{ (R_1, R_2) : \right. \\ \left. R_1 + R_2 \leq \frac{1}{2} \log \left(1 + \frac{P_1 + P_2 + 2\sqrt{P_1 P_2}}{N} \right) \right\}$$



Partial Cooperation: Conferencing Encoders

Theorem 7: (Bross/Lapidoth/Wigger'08)

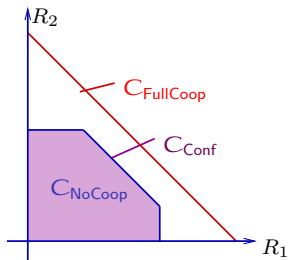
$$C_{\text{Conf}} = \bigcup_{\substack{\rho_1, \rho_2 \\ \in [0,1]}} \left\{ (R_1, R_2) : \begin{array}{l} R_1 \leq \frac{1}{2} \log \left(1 + \frac{P_1(1-\rho_1^2)}{N} \right) + C_{12} \\ R_2 \leq \frac{1}{2} \log \left(1 + \frac{P_2(1-\rho_2^2)}{N} \right) + C_{21} \\ R_1 + R_2 \leq \frac{1}{2} \log \left(1 + \frac{P_1(1-\rho_1^2) + P_2(1-\rho_2^2)}{N} \right) + C_{12} + C_{21} \\ R_1 + R_2 \leq \frac{1}{2} \log \left(1 + \frac{P_1 + P_2 + 2\rho_1\rho_2\sqrt{P_1P_2}}{N} \right) \end{array} \right\}$$

Partial Cooperation: Conferencing Encoders

Theorem 7: (Bross/Lapidoth/Wigger'08)

$$C_{\text{Conf}} = \bigcup_{\substack{\rho_1, \rho_2 \\ \in [0,1]}} \left\{ (R_1, R_2) : \begin{cases} R_1 \leq \frac{1}{2} \log \left(1 + \frac{P_1(1-\rho_1^2)}{N} \right) + C_{12} \\ R_2 \leq \frac{1}{2} \log \left(1 + \frac{P_2(1-\rho_2^2)}{N} \right) + C_{21} \\ R_1 + R_2 \leq \frac{1}{2} \log \left(1 + \frac{P_1(1-\rho_1^2) + P_2(1-\rho_2^2)}{N} \right) + C_{12} + C_{21} \\ R_1 + R_2 \leq \frac{1}{2} \log \left(1 + \frac{P_1 + P_2 + 2\rho_1\rho_2\sqrt{P_1P_2}}{N} \right) \end{cases} \right\}$$

$$C_{12} = C_{21} = 0$$

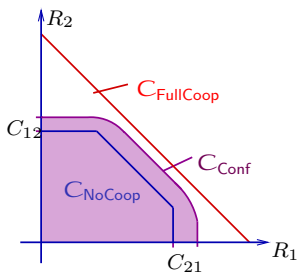


Partial Cooperation: Conferencing Encoders

Theorem 7: (Bross/Lapidoth/Wigger'08)

$$C_{\text{Conf}} = \bigcup_{\substack{\rho_1, \rho_2 \\ \in [0,1]}} \left\{ (R_1, R_2) : \begin{cases} R_1 \leq \frac{1}{2} \log \left(1 + \frac{P_1(1-\rho_1^2)}{N} \right) + C_{12} \\ R_2 \leq \frac{1}{2} \log \left(1 + \frac{P_2(1-\rho_2^2)}{N} \right) + C_{21} \\ R_1 + R_2 \leq \frac{1}{2} \log \left(1 + \frac{P_1(1-\rho_1^2) + P_2(1-\rho_2^2)}{N} \right) + C_{12} + C_{21} \\ R_1 + R_2 \leq \frac{1}{2} \log \left(1 + \frac{P_1 + P_2 + 2\rho_1\rho_2\sqrt{P_1P_2}}{N} \right) \end{cases} \right\}$$

$C_{12}, C_{21} \neq 0$,
but “small”



If $C_{12} > 0$ or $C_{21} > 0$

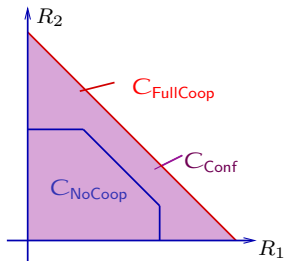
$$C_{\text{NoCoop}} \subsetneq C_{\text{Conf}}$$

Partial Cooperation: Conferencing Encoders

Theorem 7: (Bross/Lapidoth/Wigger'08)

$$C_{\text{Conf}} = \bigcup_{\substack{\rho_1, \rho_2 \\ \in [0,1]}} \left\{ (R_1, R_2) : \begin{cases} R_1 \leq \frac{1}{2} \log \left(1 + \frac{P_1(1-\rho_1^2)}{N} \right) + C_{12} \\ R_2 \leq \frac{1}{2} \log \left(1 + \frac{P_2(1-\rho_2^2)}{N} \right) + C_{21} \\ R_1 + R_2 \leq \frac{1}{2} \log \left(1 + \frac{P_1(1-\rho_1^2) + P_2(1-\rho_2^2)}{N} \right) + C_{12} + C_{21} \\ R_1 + R_2 \leq \frac{1}{2} \log \left(1 + \frac{P_1 + P_2 + 2\rho_1\rho_2\sqrt{P_1P_2}}{N} \right) \end{cases} \right\}$$

$$C_{12}, C_{21} \geq \frac{1}{2} \log \left(1 + \frac{P_1 + P_2 + 2\sqrt{P_1P_2}}{N} \right)$$



Multi-Antenna Extension

- ▶ Transmitters/receiver have multiple antennas

- ▶ $\mathbf{Y}_t = \mathbf{H}_1 \mathbf{X}_{1,t} + \mathbf{H}_2 \mathbf{X}_{2,t} + \mathbf{Z}_t;$ $\{\mathbf{Z}_t\}$ IID $\sim \mathcal{N}(\mathbf{0}, \mathbf{I})$

Theorem 8: (Wigger&Kramer'09)

Capacity region for MIMO extension:

$$\bigcup_{\substack{A_1, A_2, B_1, B_2: \\ \text{tr}(A_1 A_1^T + B_1 B_1^T) \leq P_1 \\ \text{tr}(A_2 A_2^T + B_2 B_2^T) \leq P_2}} \left\{ (R_1, R_2) : \right.$$

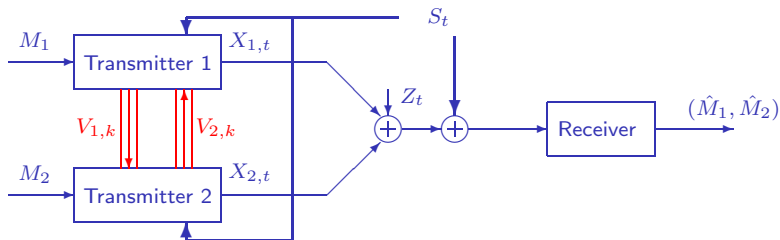
$$\left. \begin{aligned} R_1 &\leq \frac{1}{2} \log (\det (\mathbf{I} + \mathbf{H}_1 \mathbf{A}_1 \mathbf{A}_1^T \mathbf{H}_1^T)) + C_{12} \\ R_2 &\leq \frac{1}{2} \log (\det (\mathbf{I} + \mathbf{H}_2 \mathbf{A}_2 \mathbf{A}_2^T \mathbf{H}_2^T)) + C_{21} \\ R_1 + R_2 &\leq \frac{1}{2} \log (\det (\mathbf{I} + \mathbf{H}_1 \mathbf{A}_1 \mathbf{A}_1^T \mathbf{H}_1^T + \mathbf{H}_2 \mathbf{A}_2 \mathbf{A}_2^T \mathbf{H}_2^T)) \\ &\quad + C_{12} + C_{21} \\ R_1 + R_2 &\leq \frac{1}{2} \log (\det (\mathbf{I} + \mathbf{H}_1 \mathbf{A}_1 \mathbf{A}_1^T \mathbf{H}_1^T + \mathbf{H}_1 \mathbf{B}_1 \mathbf{B}_1^T \mathbf{H}_1^T \\ &\quad + \mathbf{H}_2 \mathbf{A}_2 \mathbf{A}_2^T \mathbf{H}_2^T + \mathbf{H}_2 \mathbf{B}_2 \mathbf{B}_2^T \mathbf{H}_2^T \\ &\quad + \mathbf{H}_1 \mathbf{B}_1 \mathbf{B}_2^T \mathbf{H}_2^T + \mathbf{H}_2 \mathbf{B}_2 \mathbf{B}_1^T \mathbf{H}_1^T)) \end{aligned} \right\}$$

Capacity Achieving Scheme (inspired by Willems'83)

- ▶ Transmitters split messages: $M_1 = (M_{1,c}, M_{1,p})$ and $M_2 = (M_{2,c}, M_{2,p})$
- ▶ Conference $M_{1,c}$ and $M_{2,c} \Rightarrow$ **Common Message** $(M_{1,c}, M_{2,c})$
- ▶ Rate of $M_{1,c} < C_{12}$ and rate of $M_{2,c} < C_{21}$
- ▶ Superposition $M_{1,p}$ or $M_{2,p}$ on top of $(M_{1,c}, M_{2,c})$

- ▶ MIMO: conferenced bits describe common beamforming direction

Dirty-Paper Extension



- ▶ 2 Settings: Transmitters can learn S^n **before** or **after** the conference

Theorem 9: (Bross/Lapidoth/Wigger'08)

$$C_{\text{Int,before}} = C_{\text{Int,after}} = C_{\text{Conf}}$$

Part D2: A Proof

Converse (Outer Bound on Capacity) for the
Gaussian MAC with Conferencing Encoders

Converse for Original Setup (Outer Bound)

- ▶ Step 1: Willems's outer bound *with power constraints*:

$$\mathcal{C}_{\text{Conf}} \subseteq \bigcup_{\substack{X_1-U-X_2 \\ \mathbb{E}[X_1^2] \leq P_1, \mathbb{E}[X_2^2] \leq P_2}} \mathcal{R}_{X_1, U, X_2}, \quad (1)$$

where

$$\mathcal{R}_{X_1, U, X_2} \triangleq \left\{ (R_1, R_2) : \begin{array}{l} R_1 \leq I(X_1; Y | X_2 U) + C_{12}, \\ R_2 \leq I(X_2; Y | X_1 U) + C_{21}, \\ R_1 + R_2 \leq I(X_1 X_2; Y | U) + C_{12} + C_{21}, \\ R_1 + R_2 \leq I(X_1 X_2; Y) \end{array} \right\}$$

- ▶ Step 2: In (1) suffices to take *Gaussian* Markov triples $X_1^{\mathcal{G}} - U^{\mathcal{G}} - X_2^{\mathcal{G}}$
- ▶ Step 3: Evaluate $\mathcal{R}_{X_1, U, X_2} \forall$ Gaussian Markov triples $X_1^{\mathcal{G}} - U^{\mathcal{G}} - X_2^{\mathcal{G}}$

Optimization subject to Markov Constraints

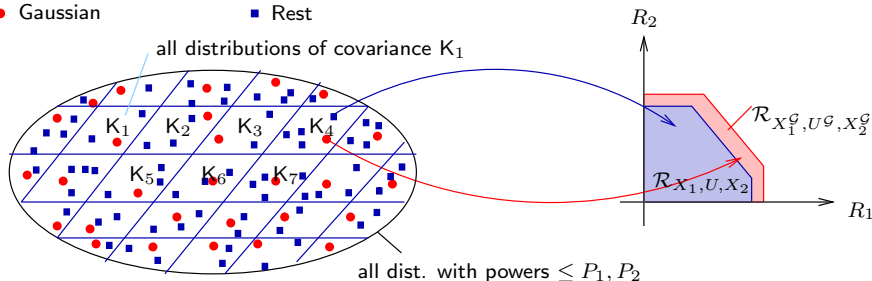
Step 2: Gaussians $X_1^G - U^G - X_2^G$ are optimal for $\mathcal{R}_{X_1, U, X_2}$:

$$\bigcup_{\substack{X_1 - U - X_2 \\ \mathbb{E}[X_1^2] \leq P_1, \mathbb{E}[X_2^2] \leq P_2}} \mathcal{R}_{X_1, U, X_2} = \bigcup_{\substack{X_1^G - U^G - X_2^G \\ \mathbb{E}[(X_1^G)^2] \leq P_1, \mathbb{E}[(X_2^G)^2] \leq P_2}} \mathcal{R}_{X_1^G, U^G, X_2^G}$$

If there were **no Markov condition**:

• Gaussian

■ Rest



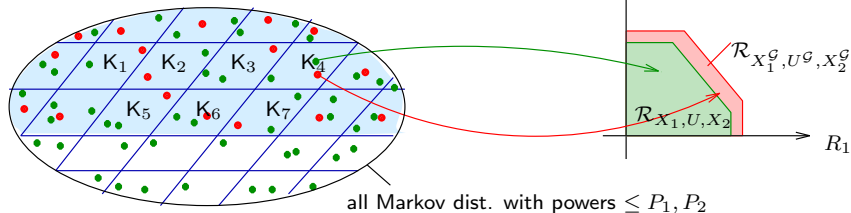
Optimization subject to Markov Constraints

Step 2: Gaussians $X_1^G - U^G - X_2^G$ are optimal for $\mathcal{R}_{X_1, U, X_2}$:

$$\bigcup_{\substack{X_1 - U - X_2 \\ \mathbb{E}[X_1^2] \leq P_1, \mathbb{E}[X_2^2] \leq P_2}} \mathcal{R}_{X_1, U, X_2} = \bigcup_{\substack{X_1^G - U^G - X_2^G \\ \mathbb{E}[(X_1^G)^2] \leq P_1, \mathbb{E}[(X_2^G)^2] \leq P_2}} \mathcal{R}_{X_1^G, U^G, X_2^G}$$

First try **with** Markov condition \rightarrow same as without Markov condition

- Gaussian Markov
- Non-Gaussian Markov



Problem: $\forall K \succeq 0$ there is a Markov triple but not necess. a Gaussian Markov!

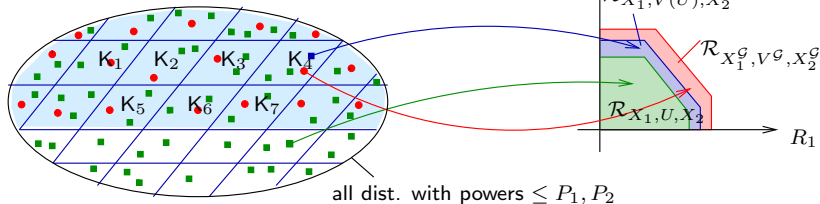
Optimization subject to Markov Constraints

Step 2: Gaussians $X_1^G - U^G - X_2^G$ are optimal for $\mathcal{R}_{X_1, U, X_2}$:

$$\bigcup_{\substack{X_1 - U - X_2 \\ \mathbb{E}[X_1^2] \leq P_1, \mathbb{E}[X_2^2] \leq P_2}} \mathcal{R}_{X_1, U, X_2} = \bigcup_{\substack{X_1^G - U^G - X_2^G \\ \mathbb{E}[(X_1^G)^2] \leq P_1, \mathbb{E}[(X_2^G)^2] \leq P_2}} \mathcal{R}_{X_1^G, U^G, X_2^G}$$

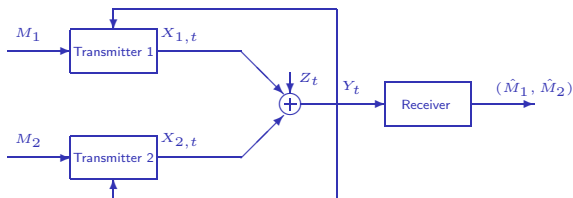
Trick: Consider X_1, V, X_2 , where $V = \mathbb{E}[X_1|U] - \mathbb{E}[X_1]$

- Gaussian Markov
- Non-Gaussian Markov
- Rest



Because for covariance of X_1, V, X_2 there is a Gaussian Markov triple □

New Tool also Applies for Cover-Leung Region



Achievable region for Gaussian MAC with perfect partial feedback

$$\mathcal{C}_{\text{PerfectFB}} \supseteq \mathcal{C}_{\text{PerfectPartialFB}} \supseteq \mathcal{R}_{\text{CL}}$$

$$\mathcal{R}_{\text{CL}} \triangleq \bigcup_{\substack{X_1 - U - X_2 \\ \mathbb{E}[X_1^2] \leq P_1, \\ \mathbb{E}[X_2^2] \leq P_2}} \left\{ (R_1, R_2) : \begin{array}{l} R_1 \leq I(X_1; Y | X_2 U) \\ R_2 \leq I(X_2; Y | X_1 U) \\ R_1 + R_2 \leq I(X_1 X_2; Y) \end{array} \right\}$$

Suffices to consider Gaussian Markov triples $X_1^G - U^G - X_2^G!$

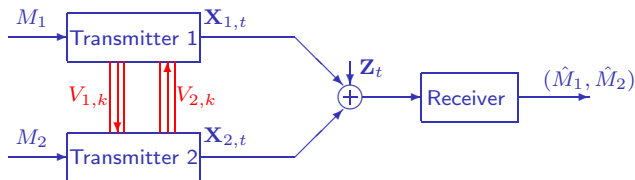
New Tool Applies for even More Settings

In expressions for capacity regions of :

- ▶ Two-users MAC with a common and two private messages (Slepian&Wolf'73)
- ▶ Interference channels with partial transmitter cooperation (Maric/Yates/Kramer'07)
- ▶ Compound MAC with conferencing encoders (Maric/Yates/Kramer'08)

it suffices to consider Gaussian Markov triples!

New Tool Extends to Vector-Case



Capacity of Gaussian MIMO MAC with Conferencing Encoders

$$C_{\text{Conf}} = \mathcal{R}_{\text{Conf}}$$

$$\mathcal{R}_{\text{Conf}} \triangleq \bigcup_{\substack{\mathbf{X}_1 - \mathbf{U} - \mathbf{X}_2 \\ \text{tr}(\mathbf{K}_{\mathbf{X}_1}) \leq P_1, \\ \text{tr}(\mathbf{K}_{\mathbf{X}_2}) \leq P_2}} \left\{ (R_1, R_2) : \begin{cases} R_1 & \leq I(\mathbf{X}_1; \mathbf{Y} | \mathbf{X}_2 \mathbf{U}) + C_{12}, \\ R_2 & \leq I(\mathbf{X}_2; \mathbf{Y} | \mathbf{X}_1 \mathbf{U}) + C_{21}, \\ R_1 + R_2 & \leq I(\mathbf{X}_1 \mathbf{X}_2; \mathbf{Y} | \mathbf{U}) + C_{12} + C_{21}, \\ R_1 + R_2 & \leq I(\mathbf{X}_1 \mathbf{X}_2; \mathbf{Y}) \end{cases} \right\}$$

Suffices to consider Gaussian Markov triples $\mathbf{X}_1^{\mathcal{G}} - \mathbf{U}^{\mathcal{G}} - \mathbf{X}_2^{\mathcal{G}}!$

New Tool Extends to Multiple Markov Chains (Wigger&Kramer'09)

Capacity of 3 Users AWGN MAC with Common Msgs (Slepian&Wolf'73)

$$C_{3\text{Users,CommonMsgs}} = \mathcal{R}_{3,\text{SW}}$$

$$\mathcal{R}_{3,\text{SW}} \triangleq \bigcup_{\substack{U_0, U_{12}, U_{13}, U_{23} \text{ indep.} \\ X_1 - (U_0, U_{12}, U_{13}) - (X_2, X_3, U_{23}) \\ X_2 - (U_0, U_{12}, U_{23}) - (X_1, X_3, U_{13}) \\ X_3 - (U_0, U_{13}, U_{23}) - (X_1, X_2, U_{12}) \\ \mathbb{E}[X_\nu^2] \leq P_\nu, \nu \in \{1, 2, 3\}}} \left\{ (R_0, R_1, R_2, R_3, R_{12}, R_{13}, R_{23}) : \right.$$

$$\begin{aligned} R_1 &\leq I(X_1; Y | X_2, X_3, U_0, U_{12}, U_{13}) \\ R_2 &\leq I(X_2; Y | X_1, X_3, U_0, U_{12}, U_{23}) \\ R_3 &\leq I(X_3; Y | X_1, X_2, U_0, U_{13}, U_{23}) \\ R_1 + R_2 &\leq I(X_1, X_2; Y | X_3, U_0, U_{12}, U_{13}, U_{23}) \\ R_1 + R_3 &\leq I(X_1, X_3; Y | X_2, U_0, U_{12}, U_{13}, U_{23}) \\ &\dots \quad \dots \\ R_0 + R_{12} + R_{13} + R_{23} + R_1 + R_2 + R_3 &\leq I(X_1, X_2, X_3; Y) \end{aligned} \left. \right\}$$

Suffices to consider Gaussians satisfying independence and Markov conditions!

Summary of Talk

Fading MIMO BC with Channel State Information @ Tx/Rxs

- ▶ Imprecisions in CSI @ tx \Rightarrow degrees of freedom collapse from 2 to $\leq \frac{4}{3}$

BC with Correlated Noises and Feedback

- ▶ 2 degrees of freedom with 1 tx-antenna and perfect fb for $\rho_z = -1$
- ▶ Noisy feedback \Rightarrow degrees of freedom collapse to 1

MAC with Feedback

- ▶ Almost noise-free feedback \approx noise-free feedback
- ▶ Even noisy feedback is always beneficial
- ▶ Answer van der Meulen's question

MAC with Conferencing Encoders

- ▶ Capacity region (also for MIMO and Costa extensions)
- ▶ Solved Optimization problem subject to Markovity conditions

Other Research Topics

- ▶ Cognitive Interference Networks, Wyner's Linear Network
"Equivalence cognition at txs and joint processing at rxs"
[Lapidoth/Shamai/Wigger ISIT'07 & ITW'07; Lapidoth/Levy/Shamai/Wigger ISIT'09]
- ▶ Relay Channels with Feedback
"With feedback, amplify&forward at relay \gg block-Markov schemes"
[Bross/Wigger, ISIT'07; Bross/Wigger, IT Jan. 2009]
- ▶ Free-Space Optical Intensity Channels
"High and low SNR asymptotics under nonnegativity, peak, and average power constraints"
[Moser/Lapidoth/Wigger, ISIT'08 and submitted to IT-Trans.]

See also: <http://people.ethz.ch/~wiggerm>