

Approximate Hypothesis Testing

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Abstract

Proof Sketch: Upper Bound

ETH

We observe samples $X = X_1, ..., X_n$ that are IID according to some unknown distribution $P \in \mathcal{H}$. How many samples are needed to guess a \hat{P} "close" to P?

We seek the approximate sample complexity

$$n_{\varepsilon,\delta}(\mathcal{H}) \coloneqq \min \left\{ n : \inf_{\hat{P}} \sup_{P \in \mathcal{H}} \mathbb{I}_{\mathbf{X} \sim \operatorname{IID} P} \left[D \Big(P, \hat{P}(\mathbf{X}) \Big) > \varepsilon \right] \le \delta \right\}$$

Main Result (Informal)

The sample complexity is at most

$$n_{\varepsilon,\delta}(\mathcal{H}) = O\left(\frac{\ln\bigl(\frac{1}{\delta}\bigr) + \ln(\mathcal{C})}{\inf_{C,C'\in\mathcal{C}} d(C,C')}\right).$$

The sample complexity is at least

First consider only distinguishing between two clusters C and C'. To that end, use a LLR test for suitable mixtures $\int_{P \in C} P \mu(P)$ vs. $\int_{P' \in C'} P' \mu'(P')$:

$$\begin{split} \sup_{C} \sup_{P \in C} \mathbb{P} \left[\hat{C}(\mathbf{X}) \neq C \right] \\ &\leq \sup_{P \in C} \mathbb{P} \left[\hat{C}(\mathbf{X}) \neq C \right] + \sup_{P' \in C'} \mathbb{P} \left[\hat{C}(\mathbf{X}) \neq C' \right] \\ &\leq 1 - \inf_{P \in C, P' \in C'} \operatorname{TV} \left(P^{\otimes n}, P'^{\otimes n} \right) \\ &\leq 1 - \inf_{P \in C, P' \in C'} h^2 \left(P^{\otimes n}, P'^{\otimes n} \right) \\ &\leq \left(1 - \inf_{P \in C, P' \in C'} h^2 (P, P') \right)^n \\ &= \left(1 - d(C, C') \right)^n \end{split}$$

$$n_{\varepsilon,\delta}(\mathcal{H}) = \Omega \Biggl(\frac{\ln \bigl(\frac{1}{\delta} \bigr)}{\inf_{C,C' \in \mathcal{C}} d(C,C')} \Biggr).$$

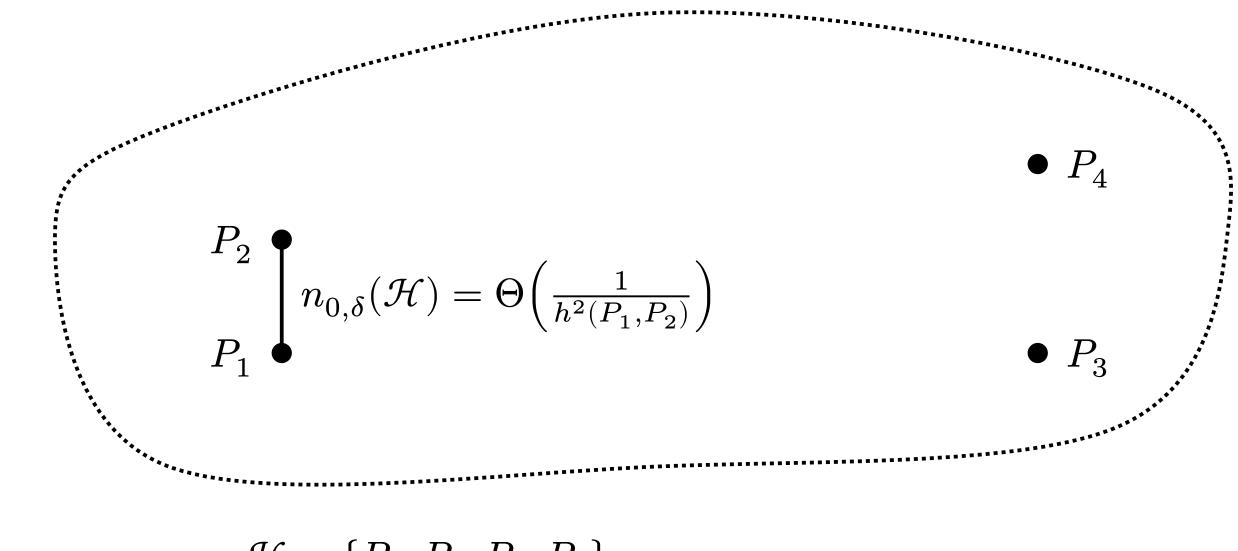
Here, \mathcal{C} is a collection of (D, ε) -dependent clusters that cover \mathcal{H} , and d(C, C') is the distance between clusters C and C' (defined below).

Classical Hypothesis Testing

Corresponds to $D(\hat{P}, P) = \mathbb{I}(\hat{P} = P)$ and $\varepsilon = 0$. Here, the sample complexity is characterized by the least squared Hellinger distance on \mathcal{H} ,

$$n_{0,\delta}(\mathcal{H}) = \Theta \Biggl(\frac{1}{\inf_{P,P' \in \mathcal{H}} h^2(P,P')} \Biggr).$$

For example:



 $\leq e^{-nd(C,C')}$

To distinguish between all clusters, apply the above test for every pair (C_i, C_j) , and take a majority vote: the cluster containing the data-generating distribution wins if it is voted for when compared to any other cluster. By the union bound, the majority vote is not won with probability at most:

$$\sup_{C} \sup_{P \in C} \mathbb{I}_{X \sim \operatorname{IID} P} \left[\hat{C}(X) \right] \leq |\mathcal{C}| e^{-n \inf_{C, C' \in \mathcal{C}} d(C, C')}.$$

Proof Sketch: Lower Bound

Observe that

- distinguishing between fixed clusters C and C' is easier than distinguishing between all clusters.
- distinguishing between fixed distributions $P \in C$ and $P' \in C'$ is easier than distinguishing between all distributions in C and C'.

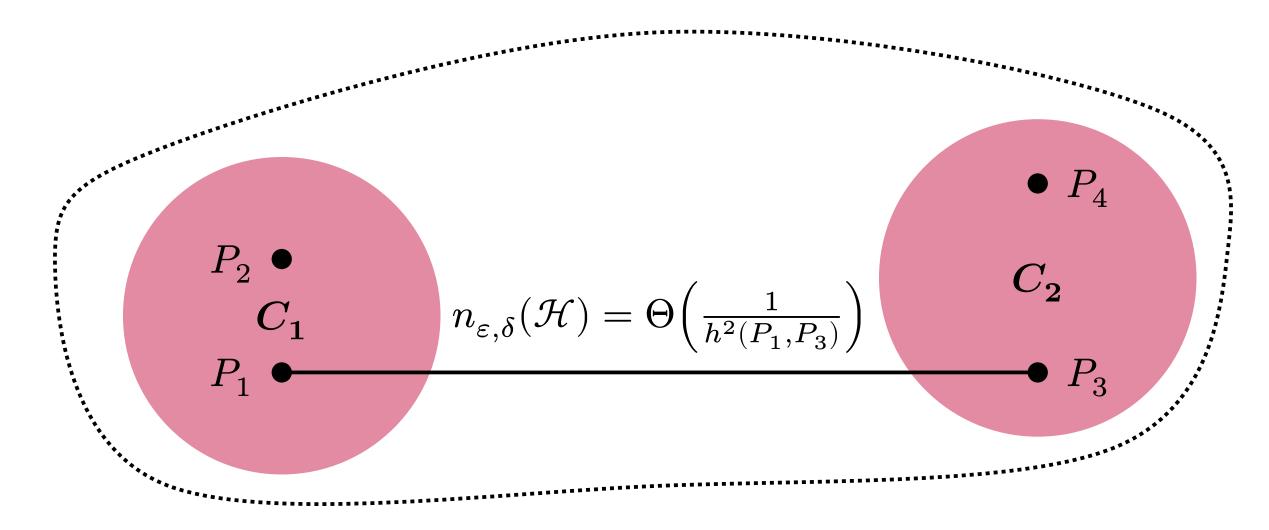
Thus, for any $C, C' \in \mathcal{C}$, and $P \in C, P' \in C'$:

 $\sup_{C} \sup_{P \in C} \mathop{\mathbb{Z}_{\sim}}_{\mathbf{X} \sim \operatorname{IID}} P^{\left[\widehat{C}(\mathbf{X})
ight]}$ $\geq \max \Big(\underset{\boldsymbol{X} \sim \operatorname{IID} P}{\mathbb{P}} \Big[\hat{P}(\boldsymbol{X}) \neq P \Big], \underset{\boldsymbol{X} \sim \operatorname{IID} P'}{\mathbb{P}} \Big[\hat{P}(\boldsymbol{X}) \neq P' \Big] \Big)$

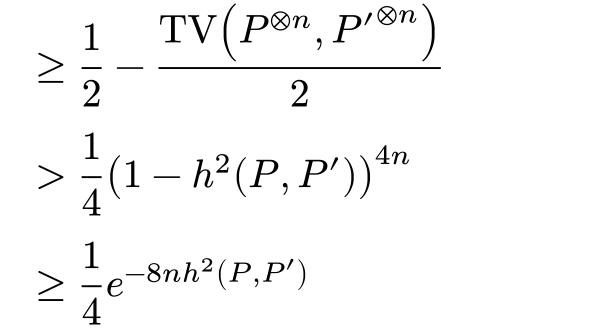
$\mathcal{H} = \{P_1, P_2, P_3, P_4\}$

Approximate Hypothesis Testing

We need no longer distinguish between P and P' that are ε -close!



 $\mathcal{H} = C_1 \cup C_2$ where $C_1 = \{P_1, P_2\}, C_2 = \{P_3, P_4\}$



Since this bound holds for any $P \in C, P' \in C'$,

$$\sup_{C} \sup_{P \in C} \mathbb{\sum}_{X \sim \operatorname{IID} P} \left[\hat{C}(X) \right] > \frac{1}{4} e^{-8nd(C,C')}.$$

And since it also holds for any $C, C' \in \mathcal{C}$,

$$\sup_{C} \sup_{P \in C} \mathbb{\sum}_{X \sim \operatorname{IID} P} \left[\hat{C}(X) \right] > \frac{1}{4} e^{-8n \inf_{C, C' \in \mathcal{C}} d(C, C')}.$$

Outlook

Key Paradigm

Approximate hypothesis testing is a flexible tool to understand data!

We plan to:

Key Insight

Cover $\mathcal H$ with clusters $\mathcal C=\{C_1,C_2,\ldots\}$ and define the distance

 $d(C,C')\coloneqq \inf_{P\in C,P'\in C'}h^2(P,P').$

Treat approximate hypothesis testing as **classical hypothesis testing on clusters** with a sample complexity that scales as $\Theta(1/\inf_{C,C'\in\mathcal{C}} d(C,C'))$.

- study the minimal cluster covering $\inf_{\mathcal{C}} |\mathcal{C}|$ as a complexity measure of the hypothesis class \mathcal{H} .
- relate the distance D(P, P') to real-world applications through $D(P, P') = |\mathbb{E}_{X \sim P}[c(X)] - \mathbb{E}_{X \sim P'}[c(X)]|^{\rho},$

where c(X) is the cost of outcome X (e.g., loss due to rising stock price).

• run simulations on synthetic and real data.

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