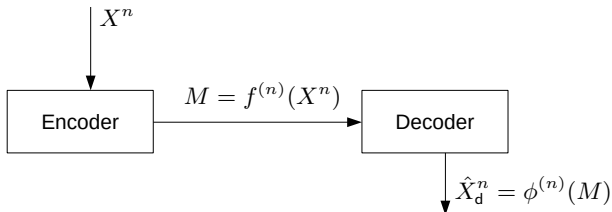


Constrained Wyner-Ziv Source Coding

Amos Lapidoth, Andreas Malär, Michèle Wigger

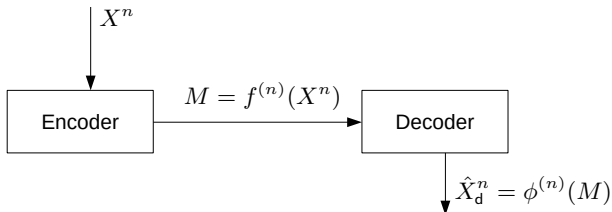
`michele.wigger@telecom-paristech.fr`

Lossless Source Coding (Perfect Reconstruction)



- ▶ Source sequence $X^n = (X_1, \dots, X_n)$ IID $\sim P_X$ over $\mathcal{X}$
- ▶ Message $M \in \{1, \dots, \lfloor 2^{nR} \rfloor\}$
- ▶ Decoder's source-reconstruction $\hat{X}_d^n \in \mathcal{X}^n$
- ▶ Error: $\hat{X}_d^n \neq X^n$

Lossless Source Coding (Perfect Reconstruction)

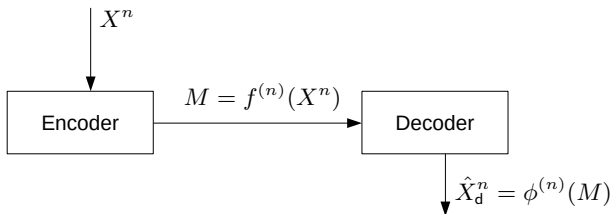


- Rate R achievable, if $\overline{\lim}_{n \rightarrow \infty} \Pr[X^n \neq \hat{X}_d^n] = 0$

Lossless Source Coding Theorem (Shannon)

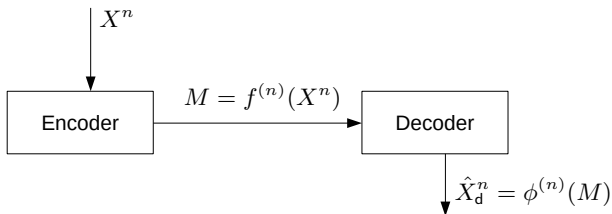
Infimum over achievable rates: $R^* = H(X)$

Lossy Source Coding (Approximate Reconstruction)



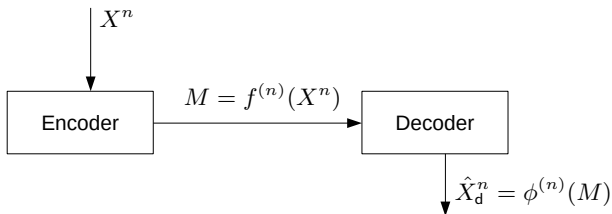
- ▶ Reconstruction alphabet $\hat{\mathcal{X}}$
- ▶ Distortion constraint $\mathbb{E} \left[d_d^{(n)}(X^n, \hat{X}_d^n) \right] \leq D_d$
- ▶ Symbol-wise distortion-function $d_d: \mathcal{X} \times \hat{\mathcal{X}} \rightarrow \mathbb{R}_0^+$
- ▶ Average per-symbol distortion: $d_d^{(n)}(X^n, \hat{X}_d^n) = \frac{1}{n} \sum_{i=1}^n d_d(X_i, \hat{X}_{d,i})$

Lossy Source Coding (Approximate Reconstruction)



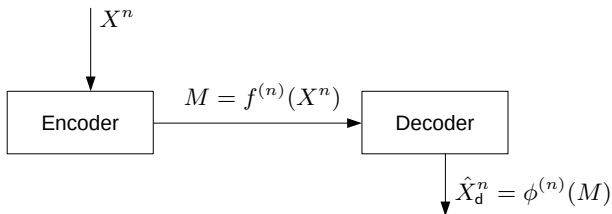
- ▶ (R, D_d) achievable if $\overline{\lim}_{n \rightarrow \infty} \mathbb{E} \left[\frac{1}{n} \sum_{i=1}^n d_d(X_i, \hat{X}_{d,i}) \right] \leq D_d$
- ▶ $R(D_d)$: infimum $R \geq 0$ such that (R, D_d) achievable

Lossy Source Coding (Approximate Reconstruction)



- ▶ (R, D_d) achievable if $\overline{\lim}_{n \rightarrow \infty} \mathbb{E} \left[\frac{1}{n} \sum_{i=1}^n d_d(X_i, \hat{X}_{d,i}) \right] \leq D_d$
- ▶ $R(D_d)$: infimum $R \geq 0$ such that (R, D_d) achievable
- ▶ Lossless problem not a special case;
- ▶ If $d_d(x, \hat{x}) = I\{x \neq \hat{x}\}$ and $D_d = 0$ (average symbol-error vs. block-error prob.)

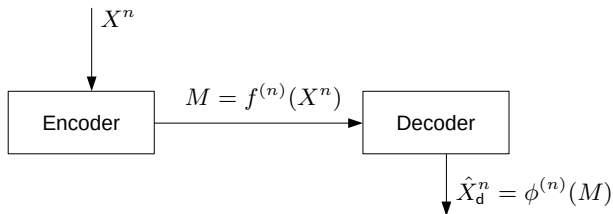
Rate-Distortion Function of Lossy Source Coding



Rate-Distortion Function

$$R(D_d) = \min_{\hat{X}_d \text{ s.t. } \mathbb{E}[d_d(X, \hat{X}_d)] \leq D_d} I(X; \hat{X}_d)$$

Rate-Distortion Function of Lossy Source Coding

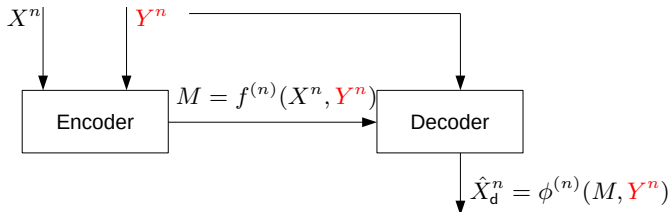


Rate-Distortion Function

$$R(D_d) = \min_{\hat{X}_d \text{ s.t. } \mathbb{E}[d_d(X, \hat{X}_d)] \leq D_d} H(X) - H(X|\hat{X}_d)$$

- Corollary: When $d_d(x, \hat{x}) = I\{x \neq \hat{x}\}$, then $R(0) = R^* = H(X)$

Lossy Source Coding with Side-Information



- Source/side-information sequence $\{(X_i, Y_i)\}$ IID $\sim P_{XY}$ over $\mathcal{X} \times \mathcal{Y}$

Rate-Distortion Function with Side-Information

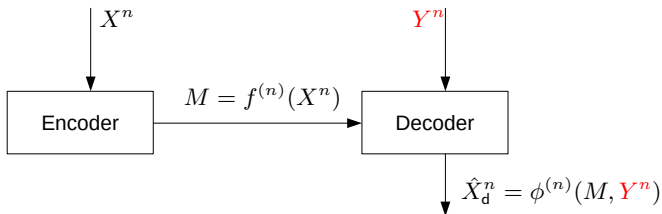
$$R_{\text{SI}}(D_d) = \min_{\hat{X}_d \text{ s.t. } \mathbb{E}[d_d(X, \hat{X}_d)] \leq D_d} I(X; \hat{X}_d | Y)$$

So far: Encoder can compute $\hat{X}_d^n$

Encoder knows decoder's reconstruction $\hat{X}_d^n$:

1. $\hat{X}_d^n$ can serve as common reference
2. If $\hat{X}_d^n$ not sufficient, encoder can re-describe X^n
→ strengthening of expected average distortion constraint

Lossy Source Coding with Side-Information @ Receiver only



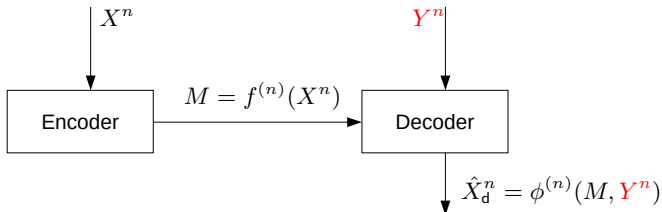
- ▶ Side-information known at decoder only!
- ▶ $R(D_d) \geq R_{WZ}(D_d) \geq R_{SI}(D_d)$

Wyner-Ziv Rate-Distortion Function

$$R_{WZ}(D_d) = \min_{\substack{Z, \hat{X}_d(Z, Y) \text{ s.t.} \\ Z \dashrightarrow X \dashrightarrow Y \\ \mathbb{E}[d_d(X, \hat{X}_d)] \leq D_d}} I(X; Z|Y)$$

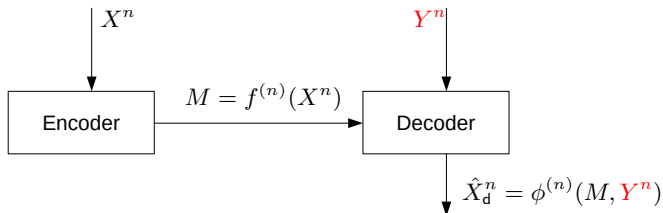
where $|\mathcal{Z}| = |\mathcal{X}| + 1$ suffices

Problem in Wyner-Ziv: Encoder cannot compute $\hat{X}_d^n$!



Encoder ignorant of $Y^n \Rightarrow$ cannot compute $\hat{X}_d^n(M, Y^n)$!

Problem in Wyner-Ziv: Encoder cannot compute $\hat{X}_d^n$!

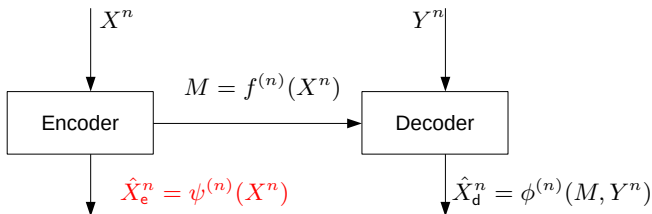


Encoder ignorant of $Y^n \Rightarrow$ cannot compute $\hat{X}_d^n(M, Y^n)$!

Wish to constrain $\hat{X}_d^n$:

- ▶ Encoder can reconstruct $\hat{X}_d^n$ **losslessly** (Steinberg'09)
- ▶ Encoder can reconstruct $\hat{X}_d^n$ **lossily** (Malär/Lapidoth/Wigger'11)

Wyner-Ziv with Common-Reconstruction (Steinberg'09)



- ▶ Same reconstruction alphabet: $\hat{X}_e^n \in \hat{\mathcal{X}}^n$
- ▶ (R, D_d) achievable if:
 1. $\overline{\lim}_{n \rightarrow \infty} \frac{1}{n} \sum_{i=1}^n \mathbb{E} \left[d_d(X_i, \hat{X}_{d,i}) \right] \leq D_d$
 2. $\overline{\lim}_{n \rightarrow \infty} \Pr \left[\hat{X}_e^n \neq \hat{X}_d^n \right] = 0$

Wyner-Ziv with Common-Reconstruction Constraint

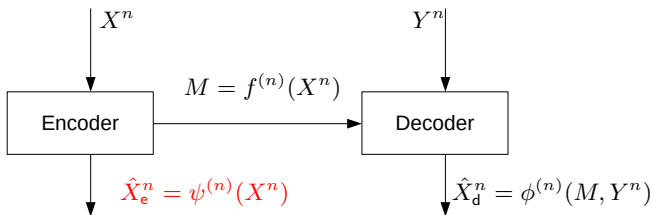
Rate-Distortion Function with Common Reconstruction (Steinberg'09)

$$R_{\text{CR}}(D_d) = \min_{\substack{Z, \hat{X}_d(Z) \text{ s.t.} \\ Z \circ - X \circ - Y \\ \mathbb{E}[d_d(X, \hat{X}_d)] \leq D_d}} I(X; Z|Y)$$

Recall Wyner-Ziv:

$$R_{\text{WZ}}(D_d) = \min_{\substack{Z, \hat{X}_d(Z, Y) \text{ s.t.} \\ Z \circ - X \circ - Y \\ \mathbb{E}[d_d(X, \hat{X}_d)] \leq D_d}} I(X; Z|Y)$$

Wyner-Ziv with **Lossy** Common-Reconstruction Constraint



- ▶ Encoder-side distortion-function $d_e: \hat{\mathcal{X}} \times \hat{\mathcal{X}} \rightarrow \mathbb{R}_0^+$
- ▶ (R, D_d, D_e) achievable if:
 1. $\overline{\lim}_{n \rightarrow \infty} \frac{1}{n} \sum_{i=1}^n \mathbb{E} \left[d_d(X_i, \hat{X}_{d,i}) \right] \leq D_d$
 2. $\overline{\lim}_{n \rightarrow \infty} \frac{1}{n} \sum_{i=1}^n \mathbb{E} \left[d_e(\hat{X}_{d,i}, \hat{X}_{e,i}) \right] \leq D_e$

Comparison with Previous Setups

Wyner-Ziv:

- ▶ Special case of our setup, e.g., when $d_d = d_e$ and $D_d \leq D_e$
 - ▶ Then choose $\hat{X}_e^n = X^n$

Steinberg:

- ▶ Not a special case of our setup $\rightarrow$ 2. not a per-symbol average
- ▶ If $d_e(\hat{x}_d, \hat{x}_e) = I\{\hat{x}_d \neq \hat{x}_e\}$ and $D_e = 0$:
block-error probability vs. average symbol-error probability

Main Result

Theorem (Rate-Distortions Function for Discrete Finite Alphabets)

$$R_{\text{lossyCR}}(D_d, D_e) = \min_{\substack{Z, \hat{X}_d(Z, Y), \hat{X}_e(Z, X) \text{ s.t.} \\ Z \dashrightarrow X \dashrightarrow Y \\ \mathbb{E}[d_d(X, \hat{X}_d)] \leq D_d \\ \mathbb{E}[d_e(\hat{X}_d, \hat{X}_e)] \leq D_e}} I(X; Z|Y)$$

where $|\mathcal{Z}| = |\mathcal{X}| + 3$ suffices

Recall Wyner-Ziv:

$$R_{\text{WZ}}(D_d) = \min_{\substack{Z, \hat{X}_d(Z, Y) \text{ s.t.} \\ Z \dashrightarrow X \dashrightarrow Y \\ \mathbb{E}[d_d(X, \hat{X}_d)] \leq D_d}} I(X; Z|Y)$$

Main Result

Theorem (Rate-Distortions Function for Discrete Finite Alphabets)

$$R_{\text{lossyCR}}(D_d, D_e) = \min_{\substack{Z, \hat{X}_d(Z, Y), \hat{X}_e(Z, X) \text{ s.t.} \\ Z \multimap X \multimap Y \\ \mathbb{E}[d_d(X, \hat{X}_d)] \leq D_d \\ \mathbb{E}[d_e(\hat{X}_d, \hat{X}_e)] \leq D_e}} I(X; Z|Y)$$

where $|\mathcal{Z}| = |\mathcal{X}| + 3$ suffices

Recall Wyner-Ziv:

$$R_{\text{WZ}}(D_d) = \min_{\substack{Z, \hat{X}_d(Z, Y) \text{ s.t.} \\ Z \multimap X \multimap Y \\ \mathbb{E}[d_d(X, \hat{X}_d)] \leq D_d}} I(X; Z|Y)$$

Corollary

When $d_e(\hat{x}_d, \hat{x}_e) = I\{\hat{x}_e \neq \hat{x}_d\}$, then $R_{\text{lossyCR}}(D_d, 0) = R_{\text{CR}}(D_d)$

Gaussian Sources and Quadratic Distortions

- ▶ $X \sim \mathcal{N}(0, \sigma_X^2)$
- ▶ $Y = X + U$, where $U \sim \mathcal{N}(0, \sigma_U^2)$ independent of X
- ▶ $d_d(x, \hat{x}_d) = (x - \hat{x}_d)^2$ and $d_e(\hat{x}_d, \hat{x}_e) = (\hat{x}_d - \hat{x}_e)^2$

Rate-Distortions Function of Gaussian Setup

Theorem (Gaussian Sources and Quadratic Distortions)

$$R_{\text{lossyCR}}(D_d, D_e) = \begin{cases} \left[\frac{1}{2} \log \left(\frac{\sigma_X^2 \sigma_U^2}{(\sigma_X^2 + \sigma_U^2) D_d} \right) \right]^+, & \text{if } \sqrt{D_e \sigma_U^2} \geq \min \left\{ D_d, \frac{\sigma_X^2 \sigma_U^2}{\sigma_X^2 + \sigma_U^2} \right\} \\ \left[\frac{1}{2} \log \left(\frac{\sigma_X^2}{\sigma_X^2 + \sigma_U^2} \frac{\sigma_U^2 + D_d - 2\sqrt{\sigma_U^2 D_e}}{D_d - D_e} \right) \right]^+, & \text{else.} \end{cases}$$

Corollary

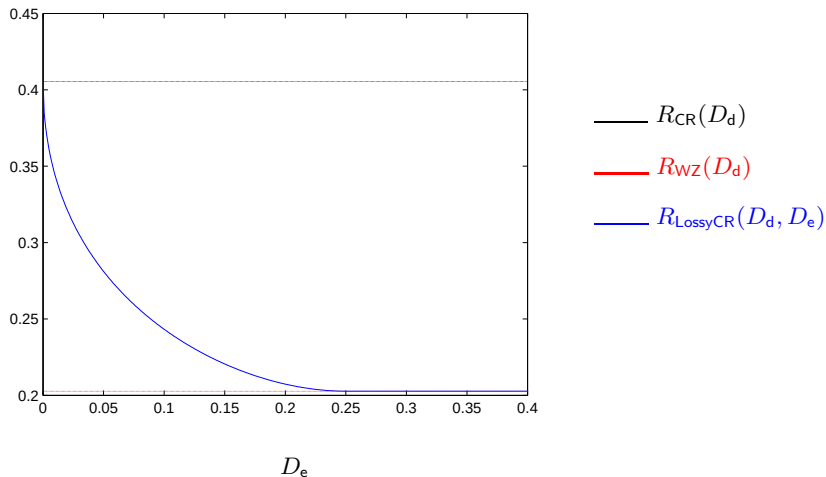
► If $\sqrt{D_e \sigma_U^2} \geq \min \left\{ D_d, \frac{\sigma_X^2 \sigma_U^2}{\sigma_X^2 + \sigma_U^2} \right\}$ or $\left(1 - \sqrt{\frac{D_e}{\sigma_U^2}}\right)^2 \sigma_X^2 \leq D_d - D_e$,

then: $R_{\text{lossyCR}} = R_{\text{WZ}} = R_{\text{SI}}$

► If $D_e = 0$, then $R_{\text{lossyCR}} = R_{\text{CR}}$

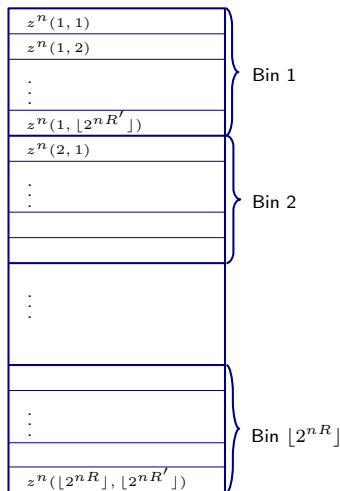
Plots for Quadratic-Gaussian Setup

$$\sigma_X^2 = 3; \quad \sigma_U^2 = 1; \quad D_d = 0.5$$



Wyner-Ziv's Scheme

Entries IID $\sim P_Z$



► Encoding:

- Choose M, K s.t.

$$(Z^n(M, K), X^n) \in \mathcal{T}_\epsilon^{(n)}(P_{ZX})$$

- Message M is bin-index!

► Decoding:

- **Binning phase:** Look for $\hat{K}$ s.t.

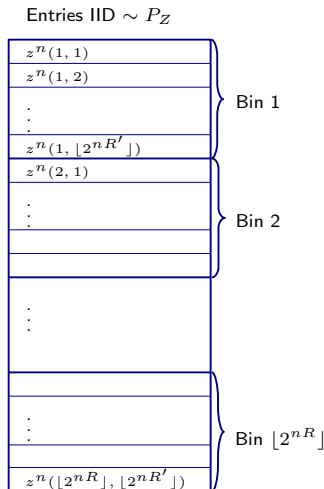
$$(Z^n(M, \hat{K}), Y^n) \in \mathcal{T}_\epsilon^{(n)}(P_{ZY})$$

- **Estimation phase:**

$$\hat{X}_{d,i} = \phi(Z_i(M, \hat{K}), Y_i)$$

With high prob: $Z^n(M, K) = Z^n(M, \hat{K})$

Steinberg's Scheme



► Encoding:

- Choose M, K s.t.

$$(Z^n(M, K), X^n) \in \mathcal{T}_\epsilon^{(n)}(P_{ZX})$$

- Message M is bin-index!

$$\hat{X}_e^n = Z^n(M, K)$$

► Decoding:

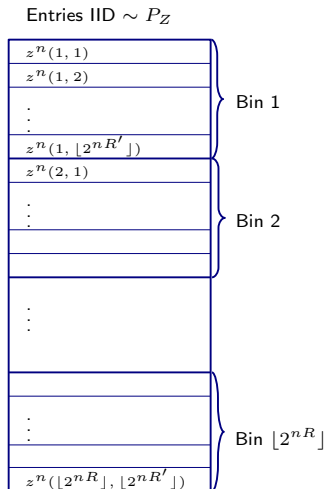
- **Binning phase:** Look for $\hat{K}$ s.t.

$$(Z^n(M, \hat{K}), Y^n) \in \mathcal{T}_\epsilon^{(n)}(P_{ZY})$$

- **"Estimation phase":** $\hat{X}_d^n = Z^n(M, \hat{K})$

Estimation phase independent of Y^n !

Our Scheme



► Encoding:

- Choose M, K s.t.

$$(Z^n(M, K), X^n) \in \mathcal{T}_\epsilon^{(n)}(P_{ZX})$$

- Message M is bin-index!

$$\hat{X}_e^n = \psi(Z_i(M, K), X_i)$$

► Decoding:

- **Binning phase:** Look for $\hat{K}$ s.t.

$$(Z^n(M, \hat{K}), Y^n) \in \mathcal{T}_\epsilon^{(n)}(P_{ZY})$$

- **Estimation phase:** $\hat{X}_{d,i} = \phi(Z_i(M, \hat{K}), Y_i)$

Estimation phase can *moderately* depend on Y^n !

Achievability in Quadratic-Gaussian Case

- ▶ Previous achievability fails (strong typicality!)
- ▶ New achievability: similar, but with coding over spheres

Converse for Discrete Case

$$\begin{aligned} \text{Converse: } R_{\text{lossyCR}}(D_d, D_e) &\geq \bar{R}(D_d, D_e) \triangleq \min_{\substack{Z, \hat{X}_d(Z, Y), \hat{X}_e(Z, X) \\ \text{s.t.: } Z \circ - X \circ - Y \\ \mathbb{E}[d_d(X, \hat{X}_d)] \leq D_d \\ \mathbb{E}[d_e(\hat{X}_d, \hat{X}_e)] \leq D_e}} I(X; Z|Y) \end{aligned}$$

a) Relax source coding problem, i.e., relax 2. distortion constraint

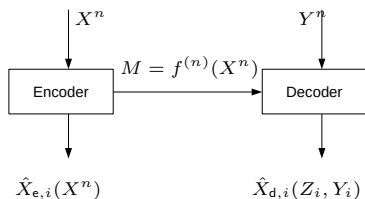
$$\text{Then: } R_{\text{lossyCR}}(D_d, D_e) \geq R_{\text{Relaxed}}(D_d, D_e)$$

b) Converse to relaxed problem:

$$R_{\text{Relaxed}}(D_d, D_e) \geq \bar{R}(D_d, D_e)$$

Converse Step a): Relax Source-Coding Problem

original problem

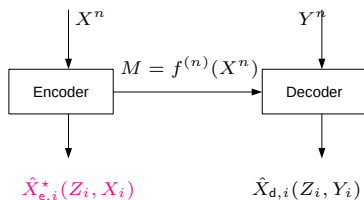


1. $\overline{\lim}_{n \rightarrow \infty} \frac{1}{n} \sum_{i=1}^n \mathbb{E} \left[d_d(X_i, \hat{X}_{d,i}) \right] \leq D_d$
2. $\overline{\lim}_{n \rightarrow \infty} \frac{1}{n} \sum_{i=1}^n \mathbb{E} \left[d_e(\hat{X}_{d,i}, \hat{X}_{e,i}) \right] \leq D_e$

► Define $Z_i \triangleq (M, Y^{i-1}, Y_{i+1}^n)$

Converse Step a): Relax Source-Coding Problem

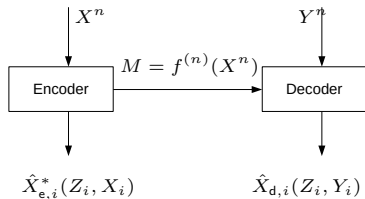
relaxed problem



1. $\overline{\lim}_{n \rightarrow \infty} \frac{1}{n} \sum_{i=1}^n \mathbb{E} [d_d(X_i, \hat{X}_{d,i})] \leq D_d$
2. $\overline{\lim}_{n \rightarrow \infty} \frac{1}{n} \sum_{i=1}^n \mathbb{E} [d_e(\hat{X}_{d,i}, \hat{X}_{e,i}^*)] \leq D_e$

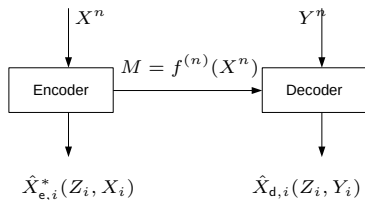
- Define $Z_i \triangleq (M, Y^{i-1}, Y_{i+1}^n)$
- Because of $X^n \text{---} (Z_i, X_i) \text{---} (Z_i, Y_i)$: new constraint 2. weaker

Converse Step b): Converse to Relaxed Problem



1. $\overline{\lim}_{n \rightarrow \infty} \frac{1}{n} \sum_{i=1}^n \mathbb{E} \left[d_d(X_i, \hat{X}_{d,i}) \right] \leq D_d$
2. $\overline{\lim}_{n \rightarrow \infty} \frac{1}{n} \sum_{i=1}^n \mathbb{E} \left[d_e(\hat{X}_{d,i}, \hat{X}_{e,i}^*) \right] \leq D_e$

Converse Step b): Converse to Relaxed Problem

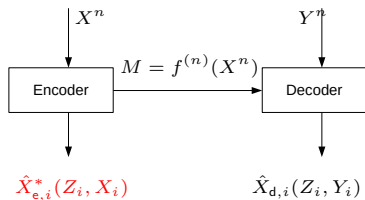


1. $\overline{\lim}_{n \rightarrow \infty} \frac{1}{n} \sum_{i=1}^n \mathbb{E} \left[d_d(X_i, \hat{X}_{d,i}) \right] \leq D_d$
2. $\overline{\lim}_{n \rightarrow \infty} \frac{1}{n} \sum_{i=1}^n \mathbb{E} \left[d_e(\hat{X}_{d,i}, \hat{X}_{e,i}^*) \right] \leq D_e$

► By definition $Z_i \triangleq (M, Y^{i-1}, Y_{i+1}^n)$: $Z_i \text{---} X_i \text{---} Y_i$

► $R_{\text{Relaxed}} \geq \frac{1}{n} H(M) \geq \dots \geq \frac{1}{n} \sum_{i=1}^n I(X_i; Z_i | Y_i)$

Converse Step b): Converse to Relaxed Problem



$$\begin{aligned}
 1. \quad & \overline{\lim}_{n \rightarrow \infty} \frac{1}{n} \sum_{i=1}^n \underbrace{\mathbb{E} \left[d_d(X_i, \hat{X}_{d,i}) \right]}_{D_{d,i}} \leq D_d \\
 2. \quad & \overline{\lim}_{n \rightarrow \infty} \frac{1}{n} \sum_{i=1}^n \underbrace{\mathbb{E} \left[d_e(\hat{X}_{d,i}, \hat{X}_{e,i}^*) \right]}_{D_{e,i}} \leq D_e
 \end{aligned}$$

► By definition $Z_i \triangleq (M, Y^{i-1}, Y_{i+1}^n)$: $Z_i \text{---} X_i \text{---} Y_i$

$$\begin{aligned}
 \text{► } R_{\text{Relaxed}} & \geq \frac{1}{n} H(M) \geq \dots \geq \frac{1}{n} \sum_{i=1}^n I(X_i; Z_i | Y_i) \geq \frac{1}{n} \sum_{i=1}^n \bar{R}(D_{d,i}, D_{e,i}) \\
 & \geq \bar{R}\left(\frac{1}{n} \sum_{i=1}^n D_{d,i}, \frac{1}{n} \sum_{i=1}^n D_{e,i}\right) \geq \bar{R}(D_d, D_e)
 \end{aligned}$$

Converse in Quadratic-Gaussian Case

- ▶ $X \sim \mathcal{N}(0, \sigma_X^2)$
- ▶ $Y = X + U$, where $U \sim \mathcal{N}(0, \sigma_U^2)$ independent of X
- ▶ $d_d(x, \hat{x}_d) = (x - \hat{x}_d)^2$ and $d_e(\hat{x}_d, \hat{x}_e) = (\hat{x}_d - \hat{x}_e)^2$

$$R_{\text{lossyCR}}(D_d, D_e) \geq \begin{cases} \left[\frac{1}{2} \log \left(\frac{\sigma_X^2 \sigma_U^2}{(\sigma_X^2 + \sigma_U^2) D_d} \right) \right]^+, & \text{if } \sqrt{D_e \sigma_U^2} \geq \min \left\{ D_d, \frac{\sigma_X^2 \sigma_U^2}{\sigma_X^2 + \sigma_U^2} \right\} \\ \left[\frac{1}{2} \log \left(\frac{\sigma_X^2}{\sigma_X^2 + \sigma_U^2} \frac{\sigma_U^2 + D_d - 2\sqrt{\sigma_U^2 D_e}}{D_d - D_e} \right) \right]^+, & \text{else.} \end{cases}$$

Converse in Quadratic-Gaussian Case, First Step

$$\begin{aligned} \text{Step 1: } R_{\text{lossyCR}}(D_d, D_e) \geq & \inf_{Z, \hat{X}_d(Z, Y), \hat{X}_e(Z, X)} I(X; Z|Y) \quad (1) \\ \text{s.t.: } & Z \text{---} \circ \text{---} X \text{---} \circ \text{---} Y \\ & \mathbb{E}[(X - \hat{X}_d)^2] \leq D_d \\ & \mathbb{E}[(\hat{X}_d - \hat{X}_e)^2] \leq D_e \end{aligned}$$

Step 2:- Evaluate RHS(1);

Converse in Quadratic-Gaussian Case, First Step

$$\text{Step 1: } R_{\text{lossyCR}}(D_d, D_e) \geq h(X|Y) - \sup_{\substack{Z, \hat{X}_d(Z, Y), \hat{X}_e(Z, X) \\ \text{s.t.: } Z \rightarrow X \rightarrow Y \\ \mathbb{E}[(X - \hat{X}_d)^2] \leq D_d \\ \mathbb{E}[(\hat{X}_d - \hat{X}_e)^2] \leq D_e}} h(X|YZ) \quad (1)$$

Step 2: Evaluate RHS(1); First Thoughts:

- ▶ Conditional Max-Entropy Theorem:
Given $K_{XYZ\hat{X}_d\hat{X}_e}$ Gaussian tuple $(Z, \hat{X}_d(Z, Y), \hat{X}_e(Z, X))$ optimizes (1)
- ▶ Not $\forall K_{XYZ\hat{X}_d\hat{X}_e}$ the Gaussian tuple is valid because $\hat{X}_d(Z, Y)$ and $\hat{X}_e(Z, X)$
- ▶ If we relax $\hat{X}_d(Z, Y)$ and $\hat{X}_e(Z, X) \Rightarrow \text{RHS}(1)=0$ (too low!)

Converse in Quadratic-Gaussian Case, Further Steps

$$\text{Step 1: } R_{\text{lossyCR}}(D_d, D_e) \geq h(X|Y) - \sup_{\substack{Z, \hat{X}_d(Z, Y), \hat{X}_e(Z, X) \\ \text{s.t.: } Z \multimap X \multimap Y \\ \mathbb{E}[(X - \hat{X}_d)^2] \leq D_d \\ \mathbb{E}[(\hat{X}_d - \hat{X}_e)^2] \leq D_e}} h(X|YZ) \quad (1)$$

Step 2: RHS(1) lower bounded by:

$$h(X|Y) - \sup_{\substack{\hat{X}_d \text{ s.t.:} \\ \mathbb{E}[(X - \hat{X}_d)^2] \leq D_d \\ \left| \mathbb{E}[(X - \hat{X}_d)U] \right| \leq \sqrt{\sigma_U^2 D_e}}} h(X - \hat{X}_d | Y - \hat{X}_d, \hat{X}_d) \quad (2)$$

Step 3: (2) maximized by jointly Gaussian $(\hat{X}_d, X, U)$ (cond. max-entropy thm)

Step 4: Evaluate (2) for jointly Gaussian $(\hat{X}_d, X, U)$

Step 2-I: Apply $\hat{X}_d(Z, Y)$ to transform Objective Function

- Because $\hat{X}_d(Z, Y)$:

$$\begin{aligned} h(X|Y, Z) &= h(X|Y, Z, \hat{X}_d) = h(X - \hat{X}_d|Y - \hat{X}_d, Z, \hat{X}_d) \\ &\leq h(X - \hat{X}_d|X - \hat{X}_d + U, \hat{X}_d) \end{aligned}$$

Step 2-I:

$$\begin{aligned} R_{\text{lossyCR}}(D_d, D_e) &\geq h(X|Y) - \sup_{\substack{Z, \hat{X}_d(Z, Y), \hat{X}_e(Z, X) \\ \text{s.t.: } Z \circ - X \circ - Y \\ \mathbb{E}[(X - \hat{X}_d)^2] \leq D_d \\ \mathbb{E}[(\hat{X}_d - \hat{X}_e)^2] \leq D_e}} h(X - \hat{X}_d|X - \hat{X}_d + U, \hat{X}_d) \end{aligned}$$

Step 2-I: Apply $\hat{X}_d(Z, Y)$ to transform Objective Function

- Because $\hat{X}_d(Z, Y)$:

$$\begin{aligned} h(X|Y, Z) &= h(X|Y, Z, \hat{X}_d) = h(X - \hat{X}_d|Y - \hat{X}_d, Z, \hat{X}_d) \\ &\leq h(X - \hat{X}_d|X - \hat{X}_d + U, \hat{X}_d) \end{aligned}$$

Step 2-I:

$$R_{\text{lossyCR}}(D_d, D_e) \geq h(X|Y) - \sup_{\substack{Z, \hat{X}_d(Z, Y), \hat{X}_e(Z, X) \\ \text{s.t.: } Z \rightarrow X \rightarrow Y \\ \mathbb{E}[(X - \hat{X}_d)^2] \leq D_d \\ \mathbb{E}[(\hat{X}_d - \hat{X}_e)^2] \leq D_e}} h(X - \hat{X}_d|X - \hat{X}_d + U, \hat{X}_d)$$

- Relax function-constraint now $\rightarrow$ Wyner-Ziv result (too loose)
- First need to use $\hat{X}_e(Z, X)$ to limit dependence of $\hat{X}_d$ on U

Step 2-II: Apply $\hat{X}_e(Z, X)$ to transform Constraints

► $Z \circ - X \circ - Y = X + U \Rightarrow (X, Z)$ ind. of U

► $\hat{X}_e(Z, X)$ & Constraint $\mathbb{E}[(\hat{X}_d - \hat{X}_e)^2] \leq D_e$:

$$\left| \mathbb{E}[\hat{X}_d \cdot U] \right| = \left| \mathbb{E}[(\hat{X}_d - \hat{X}_e)U] \right| \leq \sqrt{\sigma_U^2 D_e} \quad (3)$$

Step 2-II: relax constraints

$$R_{\text{lossyCR}}(D_d, D_e) \geq h(X|Y) - \sup_{\substack{Z, \hat{X}_d(Z, Y), \hat{X}_e(Z, X) \\ \text{s.t.: } (Z, X) \text{ ind. of } U \\ \mathbb{E}[(X - \hat{X}_d)^2] \leq D_d \\ \left| \mathbb{E}[\hat{X}_d U] \right| \leq \sqrt{\sigma_U^2 D_e}}} h(X - \hat{X}_d | X - \hat{X}_d + U, \hat{X}_d)$$

Step 2-II: Apply $\hat{X}_e(Z, X)$ to transform Constraints

► $Z \circ - X \circ - Y = X + U \Rightarrow (X, Z)$ ind. of U

► $\hat{X}_e(Z, X)$ & Constraint $\mathbb{E}[(\hat{X}_d - \hat{X}_e)^2] \leq D_e$:

$$\left| \mathbb{E}[\hat{X}_d \cdot U] \right| = \left| \mathbb{E}[(\hat{X}_d - \hat{X}_e)U] \right| \leq \sqrt{\sigma_U^2 D_e} \quad (3)$$

Step 2-II: relax constraints

$$R_{\text{lossyCR}}(D_d, D_e) \geq h(X|Y) - \sup_{\substack{Z, \hat{X}_d(Z, Y), \hat{X}_e(Z, X) \\ \text{s.t.: } (Z, X) \text{ ind. of } U \\ \mathbb{E}[(X - \hat{X}_d)^2] \leq D_d \\ \left| \mathbb{E}[\hat{X}_d U] \right| \leq \sqrt{\sigma_U^2 D_e}}} h(X - \hat{X}_d | X - \hat{X}_d + U, \hat{X}_d)$$

► Relax function constraints now

Summary

- ▶ Wyner-Ziv source coding with 2. lossy reconstruction-constraint
- ▶ Rate-distortions function (single-letter) for discrete case
- ▶ Rate-distortions function for quadratic-Gaussian case