

Testing Against Independence with Multiple Decision Centers

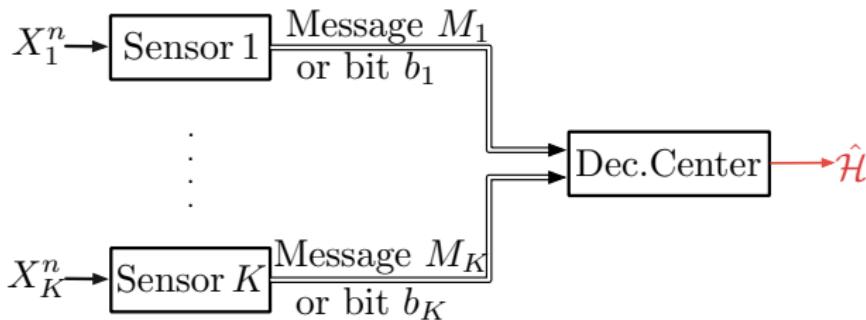
Michèle Wigger (michele.wigger@telecom-paristech.fr)

joint work with Roy Timo

SPCOM 2016, Bangalore, 13 June 2016

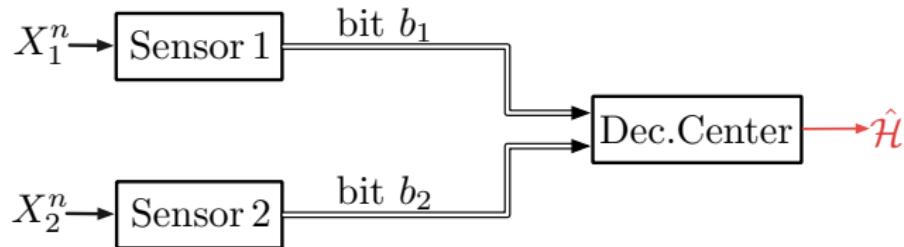
Previous setups: hypothesis testing with 1 decision center

- $(X_1^n, X_2^n, \dots, X_K^n)$ i.i.d. $\sim P_{X_1 X_2 \dots X_K}$ if $\mathcal{H} = 0$ and
 $\sim Q_{X_1 X_2 \dots X_K}$ if $\mathcal{H} = 1$



- Type-I error probability $\Pr[\hat{\mathcal{H}} = 1 | H = 0] \leq \epsilon$
- Type-II error probability $\Pr[\hat{\mathcal{H}} = 0 | H = 1] \leq e^{-\theta n}$ largest θ ?

Previous results: single-bit or zero-rate communication



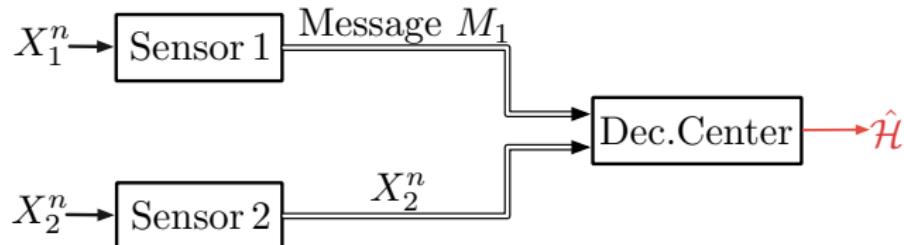
- ▶ Han['87]: Optimal error exponent:

$$\theta^* = \min_{\substack{P_{\tilde{X}_1 \tilde{X}_2} : \\ P_{\tilde{X}_1} = P_{X_1} \\ P_{\tilde{X}_2} = P_{X_2}}} D(P_{\tilde{X}_1 \tilde{X}_2} || Q_{X_1 X_2})$$

- ▶ Shalaby&Papamarcou['92]: Same with 0-rate M_1/M_2 (instead of b_1/b_2)

$$\lim_{n \rightarrow \infty} \frac{\log |\mathcal{M}_k|}{n} = 0, \quad k = 1, 2.$$

Previous results: positive-rate communication



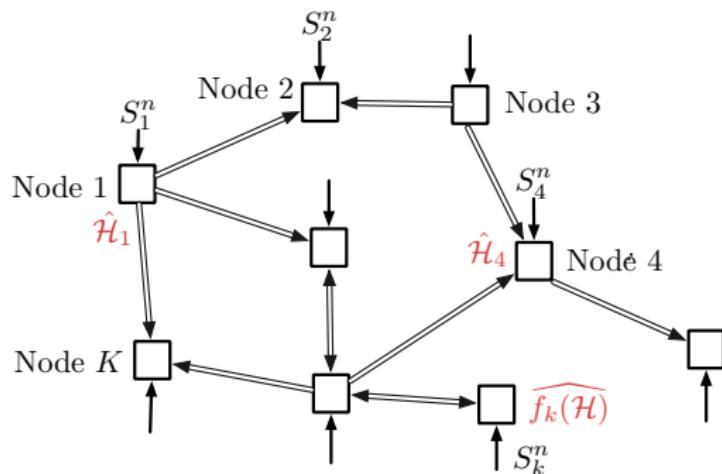
Csiszar&Ahlswede['86], Shimokawa/Han/Amari['94], Rahman&Wagner['12]:

- ▶ Test for independence $Q_{X_1 X_2} = P_{X_1} P_{X_2}$:

$$\theta^* = \max_{U \rightarrow X_1 \rightarrow X_2} I(U; X_2), \quad \text{s.t. } R \geq I(U; X_1)$$

- ▶ Extension to conditional independence test
- ▶ General achievability; converse does not match
- ▶ Interaction and secrecy: Xiang&Kim['12/'13], Zhao&Lai['14/'15], Katz/Piantanida/Couillet/Debbah['15/'16], Liao/Sankar/Tan['16]

Untouched problems in information theory



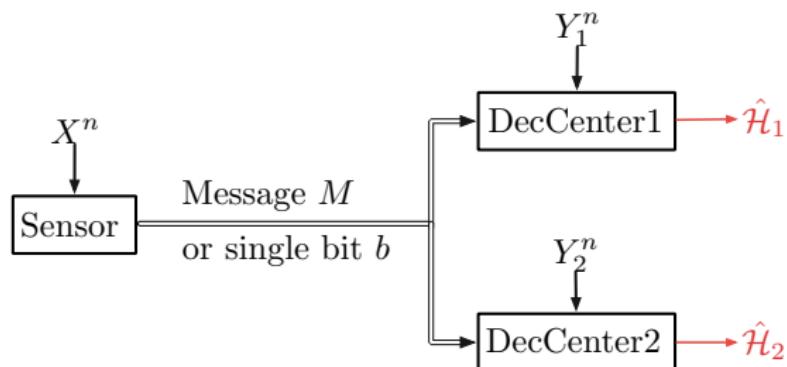
S_k^n : Node k 's measured signal

Certain nodes guess \mathcal{H}
or a function thereof

- ▶ Multiple decision centers
- ▶ Multi-hop communication

1 Sensor and 2 decision centers: the problem setup

- (X^n, Y_1^n, Y_2^n) i.i.d. $\sim P_{XY_1Y_2}$ if $\mathcal{H} = 0$ and $\sim Q_{XY_1Y_2}$ if $\mathcal{H} = 1$

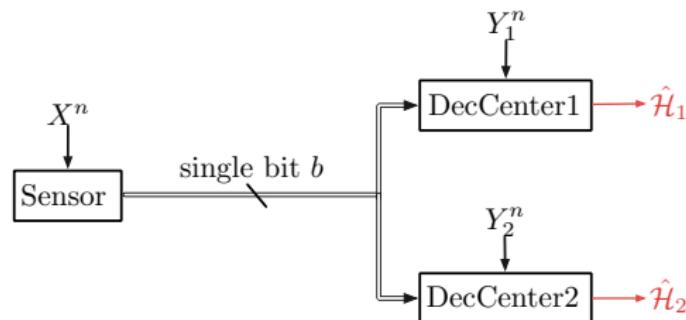


- Type-I error probabilities $\alpha_{k,n} := \Pr[\hat{H}_k = 1 | H = 0] \rightarrow 0$ as $n \rightarrow \infty$
- Type-II error probabilities $\beta_{k,n} := \Pr[\hat{H}_k = 0 | H = 1] \leq e^{-n\theta_k}$

Optimal exponents region (θ_1, θ_2) ?

Single-bit communication: the coding scheme

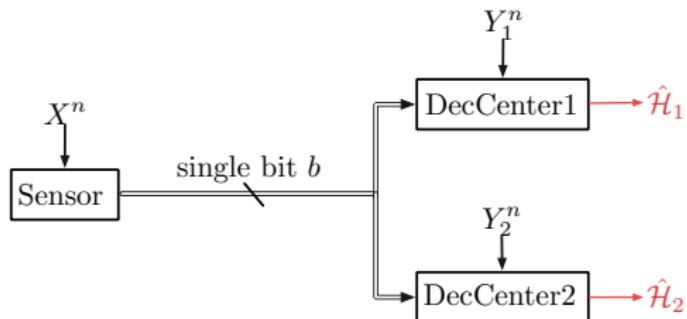
- (X^n, Y_1^n, Y_2^n) i.i.d. $\sim P_{XY_1Y_2}$ if $\mathcal{H} = 0$ and $\sim Q_{XY_1Y_2}$ if $\mathcal{H} = 1$



- Sensor: Sends $b = 1$ if X^n in typical set $\mathcal{T}_\mu(P_X)$, and $b = 0$ otherwise.
- Dec.Center k : If $(b = 1) \& (Y_k^n \in \mathcal{T}_\mu(P_Y))$, set $\hat{H}_k = 0$; otherwise, set $\hat{H}_k = 1$

Single-bit communication: optimal error-exponents

- (X^n, Y_1^n, Y_2^n) i.i.d. $\sim P_{XY_1Y_2}$ if $\mathcal{H} = 0$ and $\sim Q_{XY_1Y_2}$ if $\mathcal{H} = 1$



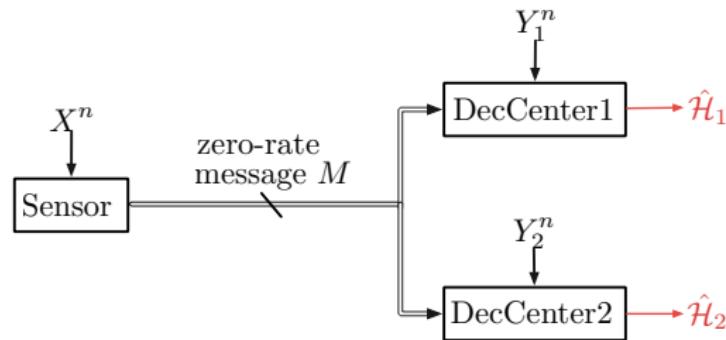
Optimal error-exponents:

$$\theta_1^* = \min_{\substack{\tilde{X}, \tilde{Y}_1: \\ P_{\tilde{X}} = P_X \\ P_{\tilde{Y}_1} = P_{Y_1}}} D(Q_{\tilde{X}\tilde{Y}_1} || P_{XY_1}); \quad \theta_2^* = \min_{\substack{\tilde{X}, \tilde{Y}_2: \\ P_{\tilde{X}} = P_X \\ P_{\tilde{Y}_2} = P_{Y_2}}} D(Q_{\tilde{X}\tilde{Y}_2} || P_{XY_2})$$

- No tradeoff because of multiple decision centers

Zero-rate communication

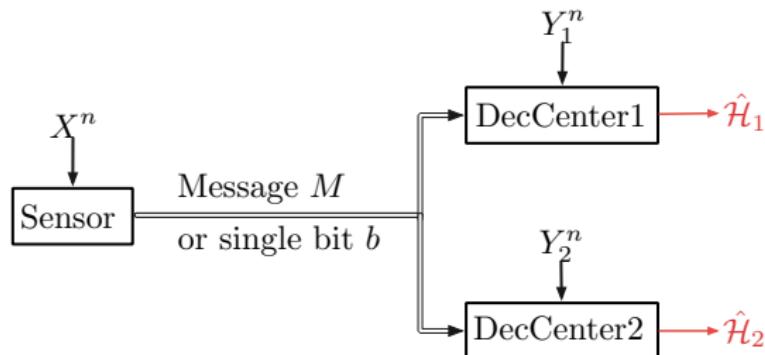
$$\lim_{n \rightarrow \infty} \frac{\log |\mathcal{M}|}{n} = 0$$



- ▶ Same exponents as with one communication bit
- ▶ No tradeoff between multiple decision centers

Positive communication rate: independence test

- (X^n, Y_1^n, Y_2^n) i.i.d. $\sim P_{XY_1Y_2}$ if $\mathcal{H} = 0$ and $\sim P_X \cdot P_{Y_1Y_2}$ if $\mathcal{H} = 1$

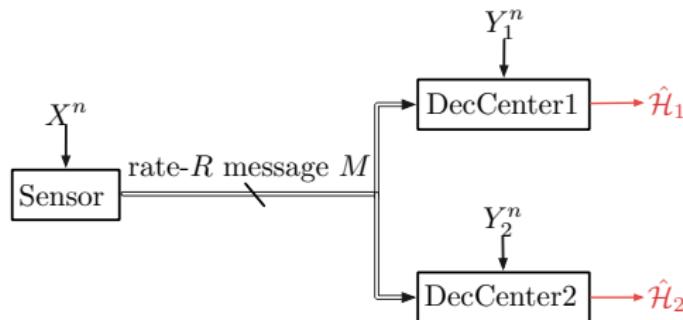


$$M \in \mathcal{M} = \{1, \dots, 2^{nR}\}$$

Rate-exponents region (R, θ_1, θ_2) ?

Test for Independence: coding scheme

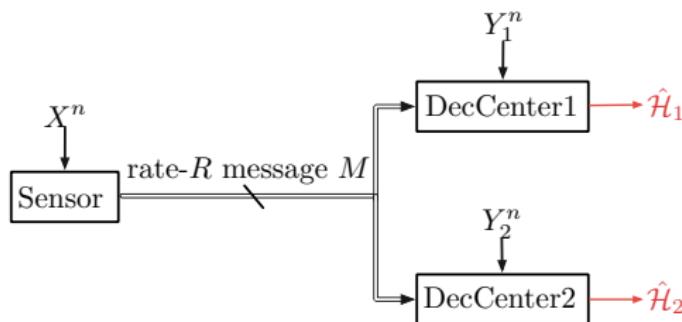
- (X^n, Y_1^n, Y_2^n) i.i.d. $\sim P_{XY_1Y_2}$ if $\mathcal{H} = 0$ and $\sim P_X \cdot P_{Y_1Y_2}$ if $\mathcal{H} = 1$



- Sensor: Quantize X^n into $U^n(m) \rightarrow$ send $M = m$
- Dec. Center k : If $(U^n(m), Y_k^n) \in \mathcal{T}_\mu(P_{UY})$, set $\hat{\mathcal{H}}_k = 0$; otherwise, set $\hat{\mathcal{H}}_k = 1$

Test for Independence: coding scheme

- (X^n, Y_1^n, Y_2^n) i.i.d. $\sim P_{XY_1Y_2}$ if $\mathcal{H} = 0$ and $\sim P_X \cdot P_{Y_1Y_2}$ if $\mathcal{H} = 1$

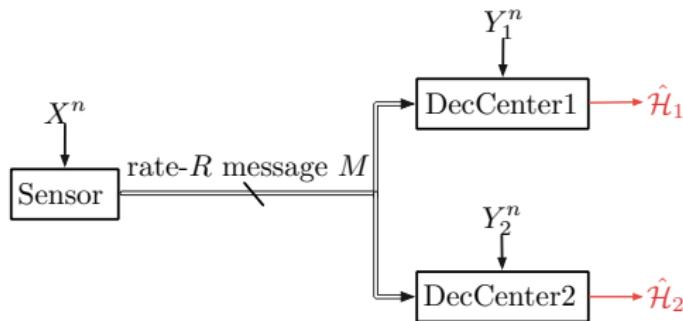


- Sensor: Quantize X^n into $U^n(m) \rightarrow$ send $M = m$
- Dec. Center k : If $(U^n(m), Y_k^n) \in \mathcal{T}_\mu(P_{UY})$, set $\hat{\mathcal{H}}_k = 0$; otherwise, set $\hat{\mathcal{H}}_k = 1$

→ Communication completely ignores receiver side-information!

Test for independence: optimal error exponents

- (X^n, Y_1^n, Y_2^n) i.i.d. $\sim P_{XY_1Y_2}$ if $\mathcal{H} = 0$ and $\sim P_X \cdot P_{Y_1Y_2}$ if $\mathcal{H} = 1$



Rate-exponents region:

$$\begin{aligned}\theta_1 &\leq I(U; Y_1) & \theta_2 &\leq I(U; Y_2) & R &> I(U; X) \\ \text{for some } U &\rightarrow X \rightarrow (Y_1, Y_2)\end{aligned}$$

- No tradeoff between decisions except for common quantisation U

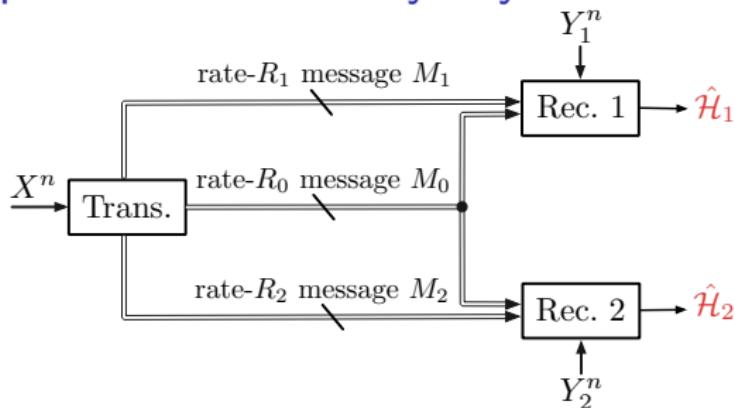
Test for independence: the converse

$$\begin{aligned} & -\frac{1}{n} \log \beta_{1,n} \\ & \leq (1-\epsilon)^{-1} \frac{1}{n} D(P_{MY_1^n|\mathcal{H}=0} \| P_{MY_1^n|\mathcal{H}=1}) \\ & = (1-\epsilon)^{-1} \frac{1}{n} I(M; Y_1^n) \\ & = (1-\epsilon)^{-1} \frac{1}{n} \sum_{t=1}^n I(M, Y_1^{t-1}; Y_{1,t}) \\ & \leq (1-\epsilon)^{-1} \sum_{t=1}^n I(M, Y_1^{t-1}, Y_2^{t-1}; Y_{1,t}) \\ & = (1-\epsilon)^{-1} I(U_n; Y_{1,n}). \end{aligned}$$

Test for independence: the converse

$$\begin{aligned} & -\frac{1}{n} \log \beta_{1,n} \\ & \leq (1-\epsilon)^{-1} \frac{1}{n} D(P_{MY_1^n|\mathcal{H}=0} \| P_{MY_1^n|\mathcal{H}=1}) & R \geq \frac{1}{n} H(M) \\ & = (1-\epsilon)^{-1} \frac{1}{n} I(M; Y_1^n) & \geq \frac{1}{n} I(M; X^n, Y_1^n, Y_2^n) \\ & = (1-\epsilon)^{-1} \frac{1}{n} \sum_{t=1}^n I(M, Y_1^{t-1}; Y_{1,t}) & \geq \frac{1}{n} \sum_{t=1}^n I(M, Y_1^{t-1}, Y_2^{t-1}; X_t) \\ & \leq (1-\epsilon)^{-1} \sum_{t=1}^n I(M, Y_1^{t-1}, Y_2^{t-1}; Y_{1,t}) & = I(U_n; X_n), \\ & = (1-\epsilon)^{-1} I(U_n; Y_{1,n}). \end{aligned}$$

Test for independence over Gray-Wynne network



Optimal rate-exponents region:

$$\theta_1 \leq I(U_0; U_1; Y_1) \quad \theta_2 \leq I(U_0; U_2; Y_2)$$

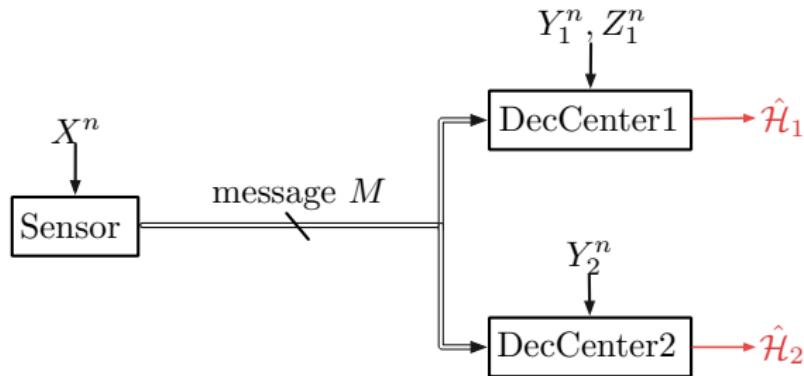
$$R_0 \geq I(U_0; X) \quad R_1 \geq I(U_1; X|U_0) \quad R_2 \geq I(U_2; X|U_0)$$

for some $(U_0, U_1, U_2) \rightarrow X \rightarrow (Y_1, Y_2)$

- ▶ No tradeoff between decisions except for common quantisation U_0
 - ▶ Receiver side-information completely ignored in communication!

An example where things change: conditional indep.

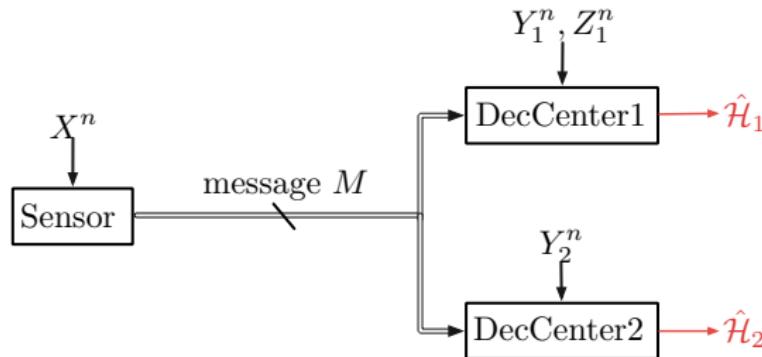
- $(X^n, Y_1^n, Y_2^n, Z_1^n)$ i.i.d. $\sim P_{XY_1Y_2Z_1}$ if $\mathcal{H} = 0$;
 $\sim P_{XZ_1} \cdot P_{Y_1|Z_1} P_{Y_2}$ if $\mathcal{H} = 1$



- Sensor: Use Heegard-Berger (binning) with side-information Z_1^n to lossily communicate X^n to both centers
- Dec. Center 1: Uses Z_1^n to recover quantization; then tests against Y_1^n

Conditional independence: Optimal Performance

- $(X^n, Y_1^n, Y_2^n, Z_1^n)$ i.i.d. $\sim P_{XY_1Y_2Z_1}$ if $\mathcal{H} = 0$;
 $\sim P_{XZ_1} \cdot P_{Y_1|Z_1} P_{Y_2}$ if $\mathcal{H} = 1$



Optimal rate-exponents region: tradeoff between centers!

$$\theta_1 \leq I(U_0, U_1; Y_1 | Z_1) \quad \theta_2 \leq I(U_0; Y_2)$$

$$R_0 \geq I(U_0; X) + I(U_1; X | U_0, Z_1) \quad \text{for some } (U_0, U_1) \rightarrow X \rightarrow (Y_1, Y_2)$$

General hypothesis tests with positive rates

- ▶ (X^n, Y_1^n, Y_2^n) i.i.d. $\sim P_{XY_1Y_2}$ if $\mathcal{H} = 0$;
 $\sim Q_{XY_1Y_2}$ if $\mathcal{H} = 1$
- ▶ Send different quantizations (U_0^n, U_1^n) and (U_0^n, U_2^n) to decision centers 1 and 2
- ▶ Adapt bin sizes to quality of $Q_{XY_1Y_2}$
→ The better Y_1 and Y_2 under Q , the less needs to be transmitted!

Communication wants to exploit side-information under $Q_{XY_1Y_2}$
Important tradeoff due to multiple decision centers

Summary

- ▶ Information-theoretic hypothesis testing with two decision centers
- ▶ Single bit or zero-rate communication:
 - no tradeoff due to multiple decision centers
- ▶ Testing for independence with positive rate:
 - some tradeoff due to multiple centers (choice of quantization);
 - optimal communication ignores receiver side-information
- ▶ Other tests at positive rates:
 - tradeoff due to multiple centers becomes important;
 - optimal communication accounts for receiver side-information