

Δ -Reliable Broadcast: A Probabilistic Measure of Broadcast Reliability

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Outline

1. The problem: scalable “reliable” broadcast
2. Reliable Broadcast specification [HT94].
3. Δ -Reliable Broadcast.
4. Reliability distribution function.
5. Examples: reliability analysis of Bimodal Multicast [BHO⁺99] and IP Multicast [DC90].

Broadcast protocols

- Best-effort: Multicast Usenet (MUSE), IP Multicast, Reliable Multicast Transfer Protocol (RMTP), etc.
- Probabilistic: Bimodal Multicast, lpbcast, etc.

What is the problem?

Traditional specification: Reliable Broadcast[HT94]

Integrity. For any message m , every process delivers m at most once, and only if m was previously broadcast by $\text{sender}(m)$.

Validity. If a correct process p broadcasts a message m , then p eventually delivers m .

Agreement. If a correct process delivers a message m , then every correct process eventually delivers m .

Informal specification: Atomicity [BHO⁺99]

A broadcast protocol provides a bimodal delivery guarantee if there is

- a high probability that a broadcast message will reach almost all processes,
- a low probability that a broadcast message will reach just a very small set of processes, and
- a vanishingly small probability that a broadcast message will reach some intermediate number of processes.

Bridging the gap: Δ -Reliable Broadcast

Let $\Delta = (\psi, \rho) \in [0, 1] \times [0, 1]$. A broadcast protocol is Δ -*Reliable* iff the following properties are *simultaneously satisfied with probability ψ* :

Integrity. For any message m , every process delivers m at most once, and only if m was previously broadcast by $sender(m)$.

Validity. If a correct process p broadcasts a message m then p eventually delivers m .

Δ -Agreement. If a correct process delivers a message m , then eventually at least a fraction ρ of correct processes deliver m .

Δ -Reliable Broadcast: ρ and ψ

$\Delta = (\psi, \rho)$ is a “reliability measure” of a given protocol.

Reliability degree ρ : the fraction of correct processes that eventually deliver a broadcast message.

Reliability probability ψ : - the probability that “enough” (correct) processes deliver a broadcast message and no fake or duplicate messages are delivered.

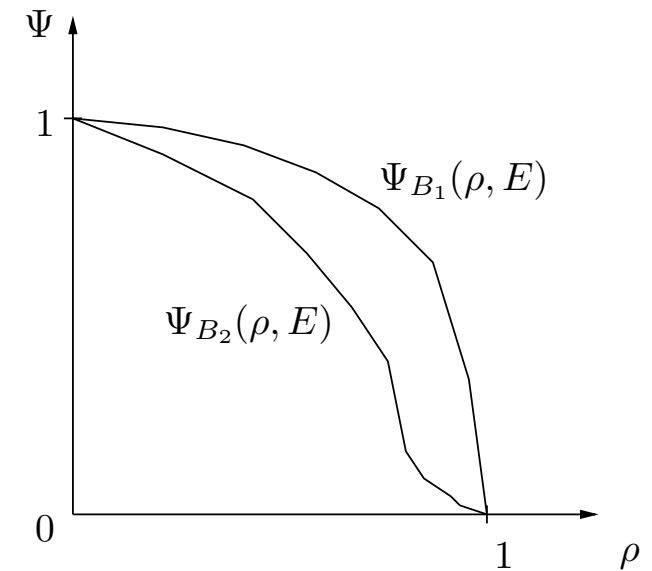
Reliability distribution function

Let $\mathbb{E}$ be an *environment space* and B be a broadcast protocol.

A function $\psi_B : [0, 1] \times \mathbb{E} \mapsto [0, 1]$ is the *reliability distribution function* of B iff

$\forall \rho \in [0, 1] \ \forall \mathcal{E} \in \mathbb{E}$:

B is Δ -Reliable with $\Delta = (\psi_B(\rho, \mathcal{E}), \rho)$.



B_1 is more reliable than B_2 in E .

Reliability distribution function: examples

- *Dreamcast* (Reliable Broadcast [HT94]) in a given $\mathcal{E} \in \mathbb{E}$:

$$\forall \rho \in [0, 1] : \psi(\rho, \mathcal{E}) = 1.$$

- *Spellcast* (does nothing):

$$\forall \rho \in]0, 1], \forall \mathcal{E} \in \mathbb{E} : \psi(\rho, \mathcal{E}) = 0.$$

Atomicity

Atomicity predicate of Bimodal Multicast: given a protocol B , $\sigma \in [0, 0.5]$ and an environment $\mathcal{E} \in \mathbb{E}$, a broadcast message reaches more than a fraction σ , but less than a fraction $1 - \sigma$ of correct processes with probability:

$$P(\sigma \leq \rho < 1 - \sigma) = \psi_B(\sigma, \mathcal{E}) - \psi_B(1 - \sigma, \mathcal{E})$$

Bimodal Multicast [BHO⁺99]

Environment:

n processes

Fanout β

Message loss probability ε

Process crash probability τ

Number of gossip rounds T

deliver_and_gossip(m , round) {^{*} Auxiliary function ^{*}}

if received_already(m) then return

bmdeliver(m)

if round= T then return

choose $S \subset \Pi$, such that $|S| = n\beta$

for each p in S send to p *gossip*(m ,round+1)

On bmcaster(m):

deliver_and_gossip(m ,0)

On receive gossip(m ,round):

deliver_and_gossip(m ,round)

IP Multicast (PIM-SM) [DC90, FHHK00]

Environment:

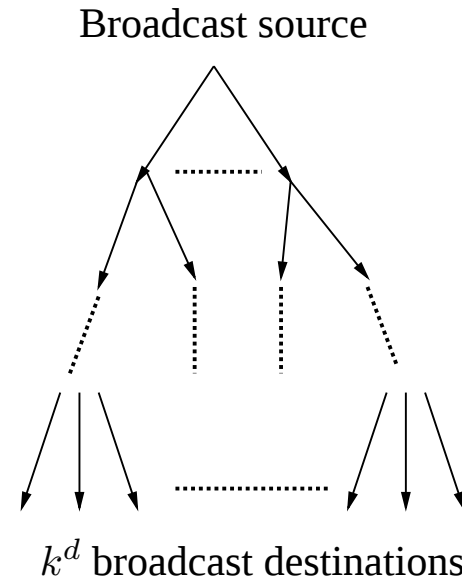
k -ary spanning tree of depth d

k^d processes

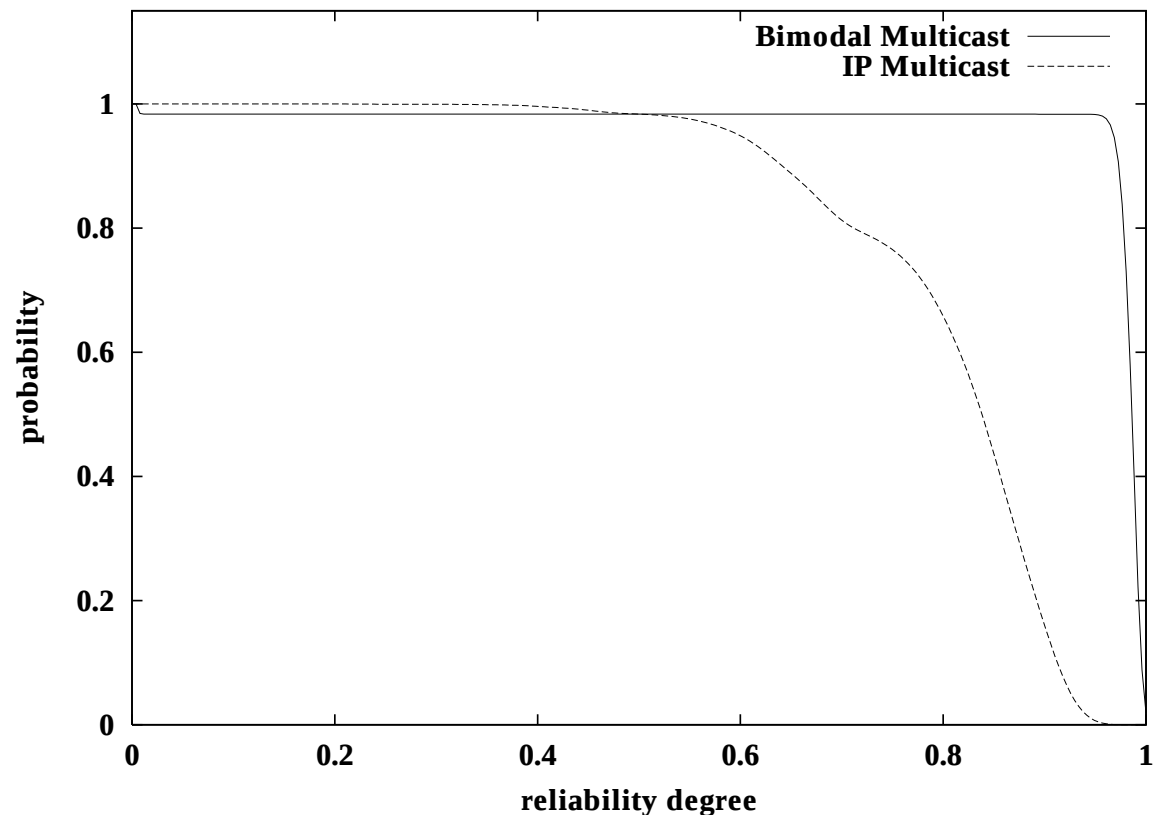
Message loss probability ε_l

Process crash probability τ

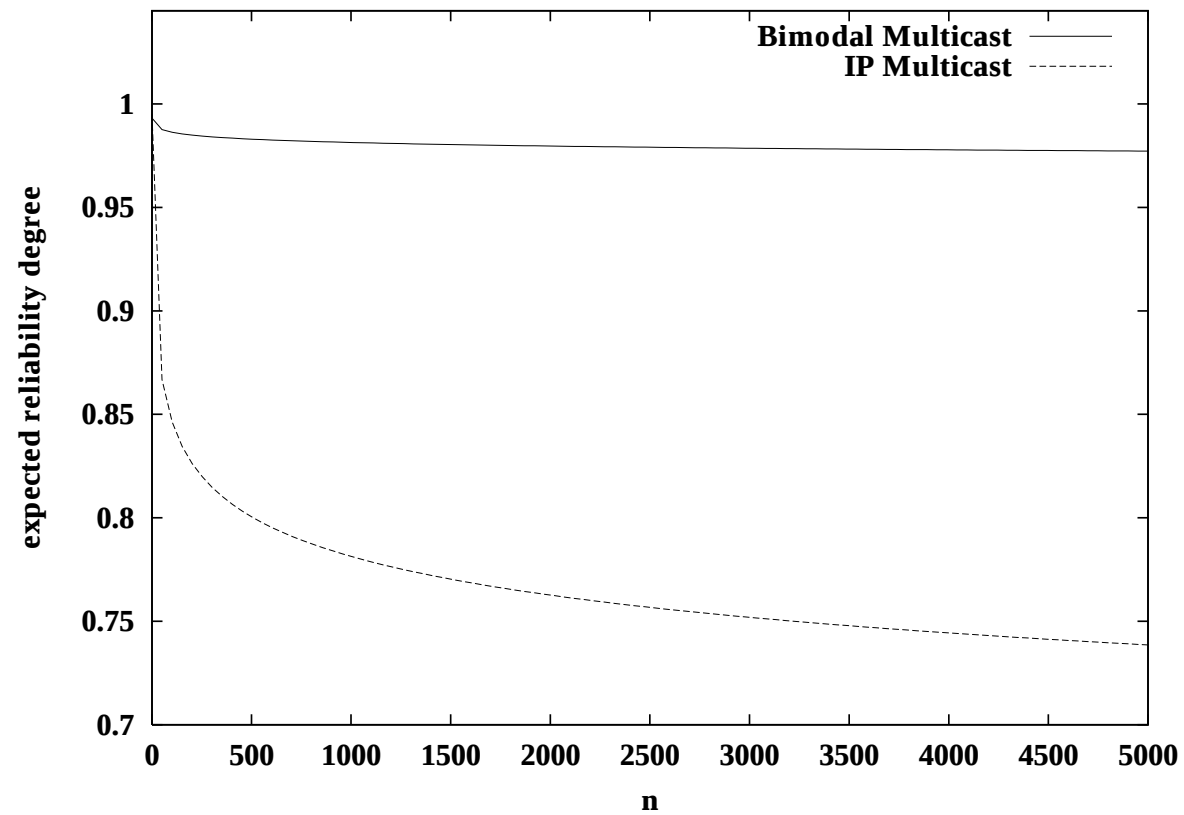
Router crash probability γ



Reliability distribution functions



Average reliability degrees



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References

- [BHO⁺99] Kenneth P. Birman, Mark Hayden, Oznur Ozkasap, Zhen Xiao, Mihai Budiu, and Yaron Minsky. Bimodal multicast. *ACM Transactions on Computer Systems*, 17(2):41–88, 1999.
- [DC90] S. Deering and D. Cheriton. Multicast Routing in Datagram Internetworks and Extended LANs. *ACM Transactions on Computer Systems*, 8(2):85–110, May 1990.
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Protocol Specification (Revised). *Internet Engineering Task Force (IETF)*, November 2000.

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