SLR210 2019 Project: Obstruction-Free Consensus and Paxos

The goal of this project is to get an initial experience in designing a fault-tolerant distributed system. Here we focus on a state-machine replicated system build atop a consensus abstraction.

1 Specification

An obstruction-free consensus (OFC) algorithm exports one operation propose(v) with an input value in a set $v \in V = \{0, 1\}$. When a process invokes propose(v), we say that the process proposes v. The operation returns either a value $v' \in V$ (in which case we say that the process decides v') or a special value *abort* $\notin V$ (in which case we say that the invocation *aborts*). A process can invoke the *propose* operation multiple times.

The following properties must be met:

- Validity: every decided value is a proposed value.
- Agreeement: no two processes decide differently.
- Obstruction-free termination:
 - If a correct process proposes, it eventually decides or aborts.
 - If a correct process decides, no correct process aborts infinitely often.
 - If there is a time after which *exactly one* correct process p proposes a value sufficiently many times, p eventually decides.

2 Concurrent environment

The goal of the project is to implement OFC for the following environment:

- We have N asynchronous processes. Every process has a distinct *identifier*. The identifiers are publicly known.
- Every two processes can communicate via a reliable asynchronous point-to-point channel.
- Up to f < N/2 of the processes are subject to crash failures: a faulty process prematurely stops taking steps of its algorithm. A process that never crashes is called correct.

3 Prerequisites

The project assumes a basic knowledge of Java. Get familiarized with the Java version of AKKA, an actor-based programming model https://akka.io/docs/. Check basic constructions in to see how to create an actor, and make the actors communicate.

Check https://github.com/remisharrock/SLR210Patterns for sample AKKA patterns which you might want to use.

4 Formalities

The project is pursued in teams of two students.

The implemented system should be provided with a short report describing how the system operates and containing correctness arguments. The team should also prepare a short presentation to be given at the end of the course.

The first project meeting on 26/04 will contain a tutorial on the AKKA programming environment and a discussion of system bootstrapping. The meeting on 17/05 will be used for discussing potential issues and problems. The final meeting on 21/06 will be used for project presentations.

5 Implementation

The implementation should extend the basic construction creating a system of a given size and ensure all-to-all connectivity. Create N actors (processes), and pass references of all N processes to each of them. ¹

Use the name **Process** for the process class. For the *Process* class, create methods for invoking the operation *propose*, processing received messages, and returning response indications.

To test the implementation and measure its performance, use the following procedure.

The main method selects f processes at random (e.g., using the shuffle method from java.collections) and sends each of them a special *crash* message. If a process receives a crash message it enters the *fault-prone* mode: for any processed *event* in the algorithm, the process decides, with a fixed probability, if it going to *crash*. If it crashes, it enters the *silent* mode, not reacting to any future event.

For every process, the main method then sends a special *launch* message. Once process *i* receives a launch message, it picks an input value, randomly chosen in $\{0, 1\}$ and invokes instances of *propose* operation with this value until a value is decided. (As a basis, one can use the OFC pseudocode to be discussed in the lecture of May 10.)

Use the LoggingAdapter class to log both the timing of the invocation and the response of every operation each process performs.

• Emulate a leader election mechanism: after a fixed timeout t_{le} , the main method randomly picks up a process that is not *fault-prone* and sends a *hold* message to every other process. After receiving a *hold* message, a process stops invoking *propose* operations.

For example, by invoking Thread.sleep(50), the main method "freezes" for 5ms.

An alternative method consists in using the scheduler. For example, the following command:

```
system.scheduler().scheduleOnce(Duration.create(50, TimeUnit.MILLISECONDS),
testActor, "foo", system.dispatcher(), null);
```

results in a message ''foo'' sent by the scheduler to testActor in 50ms.

• Perform the experiment for N = 3, 10, 100 (with f = 1, 4, 49, respectively) for different values of $t_{le} = 0.5s, 1s, 1.5s, 2s$. For each configuration, measure the time when the first process decides.

Each experiments should be repeated 5 times and the average latency should be evaluated. Build a plot realting the latency with t_{le} for different system sizes.

6 Report

Prepare a short report (up to 15 pages), preferably in English (can also be written in French if English does not feel comfortable). The report should contain:

- A high level description of the system;
- A pseudocode of the implementation;

¹Check https://gitlab.telecom-paristech.fr/petr.kuznetsov/slr210-projects.git for a example of system creation (SystemCreation). Also, the repository contains a program implementing a multi-reader multi-writer atomic register (AtomicRegister) in which every process performs a series of *read* and *write* operations on the implemented register.

- A sketch of a proof of correctness (please argue that both safety and liveness hold);
- A report on performance analysis.

The report and the code of the implementation should be uploaded to https://gitlab.telecom-paristech.fr/petr.kuznetsov/slr210-projects.git by June 14.

7 Presentation

The presentation (7 mins) should contain a brief overview of the main features of the algorithm, its correctness arguments and performance. We envision 10 minutes per team (including 3 minutes for questions), so the time bounds are strict.