

Internet Security Association and Key Management Protocol

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ISAKMP: Intentions

 Intended to support the SA management for security protocols at all layers (IPSEC, TLS, OSPF, etc.)

 Centralizes the SA management and reduces the amount of duplicated functionality



ISAKMP: Overview

- Defined in RFC 2408
- Cleanly separates the two points:
 - Security Association (SA) and key management
 - key exchange protocol (not defined here)
- Builds a framework for:
 - Transferring key and authentication data
 - Negotiation, establishment, modification and deletion of SAs



ISAKMP: Motivation

- Growing number of complex VPNs:
 - One security requirement within the VPN and many others for communications outside

- Mobility
 - Limited bandwidth
 - Need for authentication



ISAKMP: Requirements

 Linkage of authentication and key exchange within one protocol, linkage of the SA establishment with this protocol

 Usage of strong authentication with digital signatures based on PKcryptography



ISAKMP: Basic run – 2 phases

Phase 1: ISAKMP SA establishment

- Initial protocol exchange:
 - Agreement upon basic set of security attributes
 - Provides protection for subsequent exchanges
 - Indicates the auth-method and key exchange to be performed within ISAKMP
 - Can be skipped if the basic set already exists
- Authentication+key (material) exchange
- Generation of required keys

Phase 2: Protocol SA establishment



ISAKMP: Why 2 phases?

- 1. Amortization of the costs of the first phase across several second phase negotiations
- First channel provides security properties for the second phase. The security can be adjusted in the second phase.
- 3. Reduced cost of the ISAKMP management: no re-auth for each error

ISAKMP: Key exchange protocol

- Is not defined in ISAKMP, possibility: IKE
- Requirements:
 - Key generation vs. transport
 - Perfect forward secrecy (PFS)
 - Computational overhead
 - Key escrow
 - Key strength
- Negotiated and supported by ISAKMP



ISAKMP: Protection (1/2)

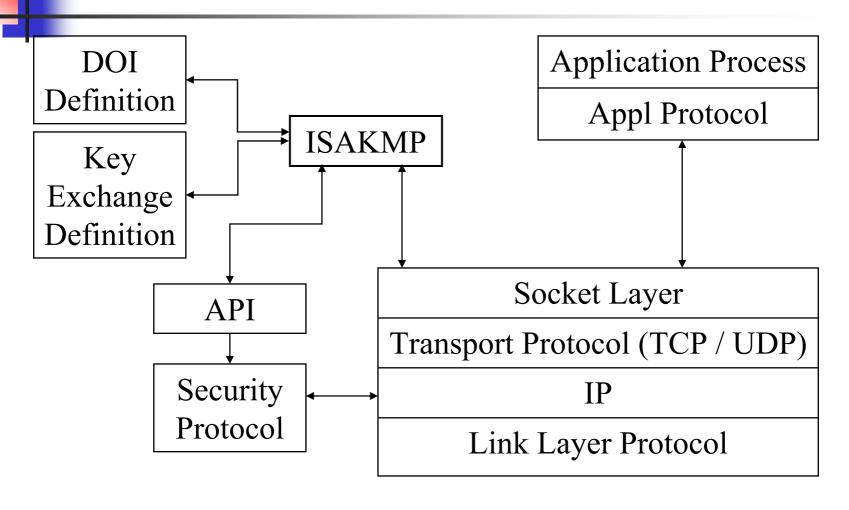
- Denial of Service
 - Cookie exchange (using only fast operations)
 - Aggressive memory management in the FSM
 - Absolute protection impossible!!!
- Connection Hijacking
 - Authentication linked with key and SA exchanges
- Multicast communications
 - Planned as a future extension



ISAKMP: Protection (2/2)

- Man-in-the-Middle
 - Exchange linking against message insertion
 - No partial SA creation
 - FSM design against reflection back to sender
 - Time variant material against replaying
 - Strong authentication
 - Detection of redirection and modification

ISAKMP: Placement





ISAKMP: Cookies

- So-called Anti-Clogging Tokens (ACT)
- Three generation rules:
 - Depend on the specific parties
 - No other entity can generate the same cookie so it will be accepted by the entity
 - 3. The generation function must be fast
- Example:

cookie := md5(IPsrc, IPdst, UDPsrc, UDPdst, RndVal, timestamp)

ISAKMP: Fixed header format

Initiator Cookie

Responder Cookie

Next Payload MjVer MnVer Exchange Type Flags

Message ID

Length

Next Payload (1 byte): The format of the next payload Major/Minor Version (4 bit): ISAKMP protocol version Exchange Type (1 byte): Type of exchange being used Flags (1 byte): Specific options set for the exchange



- Each payload type has its own header
- Header chaining with the "general header"
- Exact header formats defined in the RFC (3.4-3.16)
- Header processing rules in the RFC (5)

- Security Association
- 2. Proposal
- 3. Transform
- 4. Key Exchange
- 5. Identification
- 6. Certificate
- 7. Certificate Request
- 8. Hash
- 9. Signature
- 10. Nonce
- 11. Notification
- 12. Delete
- 13. Vendor ID
- -- Private USE (128-255)



ISAKMP: Conclusions

- Well-designed protocol aimed to the future
- Contains all the features needed for Internet's massive growth
- Flexible through negotiation
- Able to establish SAs for multiple protocols
- Can be used with arbitrary transports
- Follows the design principles of IPv6