### A journey in decision making processes

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## Who I am?

My very short resume : join S2A team in june 2023

- 1996 : Engineer degree from TP and M2 ATSI (UPSay)
- 2000 : PhD Thesis (Blind frequency and channel estimation)
- 2001 : Ass. Prof. at Digital Communications team in TP
- 2010 : Full Prof. and head of the team (until 2021)

### Some editorial activities :

- BoG Gretsi, TPC Gretsi, Icassp, Eusipco, ···
- AE/SAE in IEEE TSP, IEEE TISPN, ···

#### Some teaching activities :

- COM105, Telecom track
- SI101, MD221, M2 MICAS
- TSE202, TSE101 (next year)

#### Other TSE-related activities :

• Member of GDR Ecoinfo, GDR Internet et Société

### Outline

### Two technical topics :

- Graph Node classification
  - No Graph Neural Network (GNN)
  - Interpretable algorithm
  - Less complex algorithm (with less hyperparameters)
     H. Hafidi et al., "Graph-assisted Bayesian node classifiers", IEEE Access, 2023
- Edge caching with popular time-sensitive contents
  - No neural network (while decision making agent)
  - Low-complex interpretable probabilistic approach
     H. Tang, et al., "Cache updating strategy minimizing AoI with time-varying popularity", IEEE Information Theory Workshop, 2021

### Some perspectives :

- Wireless federative learning
- Graph based image classification
- Sustainable systems

## Topic 1 : Graph Node classification



#### Idea : homophily principle

Predict class of each unlabeled node in the graph by relying

- on nodes' features and
- on nodes' graph connections

### Examples :

- in social networks, people are more likely to connect with those who share the same areas of interest
- in research articles' database, more likely to have connections/citations between articles dealing with the same research topic

### Problem statement

Classifier based on Bayesian decision theory : Maximum A Posteriori

- $V_u$  : set of nodes involved in the classification of node u.
- $\mathcal{X}_u = \{ \mathbf{x}_u \} \cup \{ \mathbf{x}_v, v \in \mathcal{V}_u \}$ : set of features of node *u* and its "helping" nodes
- *y<sub>u</sub>* : class of node *u* (what we are looking for !)
- *D<sub>k</sub>* : probability density function of features belonging to class *k*. For any *u*,

$$D_k(\boldsymbol{x}_u) = p(\boldsymbol{x}_u|y_u = k).$$

#### Graph-Assisted Bayesian (GAB) Classifier

$$\hat{k}_u = rg\max_k P_u(k)$$

with  $P_u(k) = \Pr(y_u = k | \mathcal{X}_u, \mathcal{I}_G)$ 

### Problem solution

Derivations of  $P_u(k)$ . Bayes' rule

$$P_{u}(k) = \frac{p(\mathcal{X}_{u}|y_{u} = k, \mathcal{I}_{\mathcal{G}}) \Pr P(y_{u} = k|\mathcal{I}_{\mathcal{G}})}{P(\mathcal{X}_{u}|\mathcal{I}_{\mathcal{G}})} \propto Q_{u}(k)\pi_{k}$$

with  $\pi_k = \Pr(y_u = k | \mathcal{I}_G)$  a priori classes' probability Let  $\Delta_u$  be the diameter of the set  $\mathcal{V}_u$ .

$$Q_u(k) = D_k(\boldsymbol{x}_u) \prod_{d=1}^{\Delta_u} \prod_{v \in \mathcal{N}_u(d)} \left( \sum_{k'=1}^K r_{u,v}(k,k') D_{k'}(\boldsymbol{x}_v) \right)$$

with  $r_{u,v}(k, k') = \Pr(y_v = k' | y_u = k, \mathcal{I}_{\mathcal{G}})$  the probability to be on class k' for node v given the fact that we are in class k for node u.

#### Example

$$V_u = \{v\}$$
, known  $k_v = 1$ ,  $\pi_1 = \pi_2 = 1/2$ , and  $\Delta_u = 1$ :

$$Q_u(1) = D_1(\boldsymbol{x}_u) \frac{p}{p+q}$$
 and  $Q_u(2) = D_2(\boldsymbol{x}_u) \frac{q}{p+q}$ 

with p (resp. q) probability of intra (resp. inter)-class connection

### Assumptions

- 2 equilikely classes
  - p(k) probability that two nodes from class k are connected  $\overline{p}_{\text{arithmetic}}$  arithmetic average of  $\{p(k)\}_k$
  - *q* probability that two nodes from different classes are connected.
- Information on graph is 1-hop

We get

$$\begin{array}{c|c} r(1,2) = \frac{q}{p(1)+q} & r(2,2) = \frac{p(2)}{q+p(2)} \\ \hline r(1,1) = \frac{p(1)}{p(1)+q} & r(2,1) = \frac{q}{q+p(2)} \end{array}$$

Graph-agnostic iff r(1,2) = r(2,2) and r(1,1) = r(2,1)

Main result

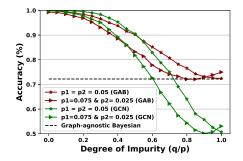
Graph-agnostic iff

• 
$$q = \sqrt{p(1)p(2)} = \overline{p}_{\text{geometric}}$$
, or

• Degree of Impurity  $= \frac{q}{\overline{\rho}_{\text{arithmetic}}} = \frac{\overline{\rho}_{\text{geometric}}}{\overline{\rho}_{\text{arithmetic}}} \le 1$ 

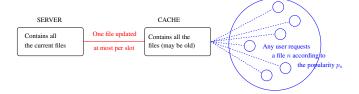
## Numerical illustrations

- 2 classes
- Gaussian distributions with different means and covariance matrices
- Number of nodes N = 5,000 and number of features F = 500
- 500 (already-labeled) nodes



- GAB more robust to Dol than GCN
- GCN becomes worse than graph-agnostic (too confident)

# Topic 2 : Edge caching



- Content n is time-sensitive (X<sub>n</sub>(t) : age in caching)
- Content n has its own popularity (p<sub>n</sub> : probability to be requested)
- Ex : newspaper website, web crawling, video last version, ...

#### Question

- Given a timeslot *t*, which item should be downloaded from the server to the cache to be as up-to-date as possible?
- Scheduling problem

 $\{u_t\}_t = f(\text{information on the system})$ 

Optimization problem

$$\arg\min_{u_1,\cdots,u_T}\sum_{n=1}^N p_n \int_0^T X_n(t) \, dt$$

s.t.  $u_t \in \{1, \dots, N\}$  for all t, and  $\sum_{t=1}^{T} \mathbf{1}\{u_t > 0\} = T$ .

#### Approaches :

- As underlying Markov chain, constrained MDP well adapted
  - Optimal random policy exists
  - Suboptimal approach but simple : Whittle's index
- Probabilistic method by re-writing the problem

**Extensions :** i) when the transmission size depends on the age, ii) when the popularity is time-varying (satisfying Markov chain)

# Approach 1 : concept of per-file update rate

- Consider  $\lambda_n$  the per-file update rate
- Actually, when T large enough,

$$\frac{1}{T}\int_0^T X_n(t)dt\approx \frac{1}{\lambda_n}$$

New optimization problem

$$\min_{\lambda_1,\ldots,\lambda_N}\sum_{n=1}^N\frac{p_n}{\lambda_n}$$

s.t. 
$$\lambda_n \geq 0$$
, and  $\lambda_1 + \cdots + \lambda_N = 1$ .

#### Main result

Problem is convex and leads to

$$\lambda_n^* = \frac{\sqrt{\rho_n}}{\sum_{m=1}^N \sqrt{\rho_m}}$$

#### Update rate of file *n* follows a square-root law wrt. its popularity

## Practical protocol

Let  $\tau_n^{\star} = 1/\lambda_n^{\star}$  be the optimal inter-update time for file *n* 

$$u_t = \arg \max_{u \in \{1, \dots, N\}} \underbrace{(X_u(t) - \tau_u^*)}_{\text{Schedule and under the Aux hand Division (S)}}$$

Schedule-ordered by Age-based Priority (SOAP)

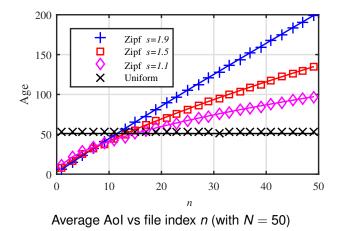
#### General context :

- r(D, X) : rank function with descriptor D and age X
- Scheduled user

$$u_t = \arg \max_{u \in \{1, \cdots, N\}} r(D_u(t), X_u(t))$$

- Many policies follow this shape
  - Round-Robin (RR),  $r(\emptyset, X_u) = X_u$
  - "Weighted Round-Robin",  $r(d_u, X_u) = d_u X_u$ Example : does it make sense to choose  $d_u = \sqrt{p_u}$ ?

## Numerical illustrations



Proposed policy reduces the average age of more popular items at the expense of less popular items

## Approach 2 : index based policy

Find a suboptimal policy based on an index :

$$u_t = \arg \max_{u \in \{1, \cdots, N\}} \mathcal{I}_u(S_u)$$

- I : it is an heuristic
- Whittle's index : methodology for exhibiting a reasonnable index in Restless Multi-Arm Bandit problem
  - N bandits/players/agents
  - At each timeslot, select one bandit (let's say  $u_t$ )
  - Its state  $s_{u_t}$  is modified according to its action, and is rewarded
  - but states of other bandits also modified and rewarded in different ways (restless)
- When non-playing bandits are frozen (no state evolution) and not rewarded : Gittins' index is optimal

$$\mathcal{IG}_u(D) = \sup_{\tau > 0} \frac{\mathbb{E}[\sum_{t=0}^{\tau-1} \gamma^t r_u(S_t) | s_0 = S]}{\sum_{t=0}^{\tau-1} \gamma^t}$$

# Whittle index (1/2)

$$\arg \max_{\{a_n(t)\}_{n,t}} \lim_{T \to \infty} \mathbb{E} \left[ \sum_{t=0}^{T-1} \gamma^t \sum_{n=1}^N r_n(s_n(t), a_n(t)) \right]$$

s.t.  $\sum_{n=1}^{N} a_n(t) = 1$  (C1) and  $a_n(t) \in \{0, 1\}$  (C2) Two modifications :

• Relaxation : C1 replaced with  $\sum_{t=0}^{\infty} \gamma^t \sum_{n=1}^{N} a_n(t) = 1/(1-\gamma)$ 

Lagrangian penalty

$$\arg \max_{\{a_n(t)\}_{n,t}} \mathcal{L}(\lambda)$$

s.t. C2 and with

$$\mathcal{L}(\lambda) = \lim_{T \to \infty} \mathbb{E}\left[\sum_{t=0}^{T-1} \gamma^t \sum_{n=1}^N r_n(s_n(t), a_n(t))\right] - \lambda\left(\sum_{n=1}^N a_n(t) - 1/(1-\gamma)\right)$$
$$= \lim_{T \to \infty} \mathbb{E}\left[\sum_{n=1}^N \left(\sum_{t=0}^{T-1} \gamma^t r_n(s_n(t), a_n(t)) - \lambda a_n(t)\right)\right]$$

# Whittle index (2/2)

- The problem is now decoupled
- For each bandit and fixed  $\lambda$ , we maximize

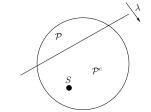
$$\mathcal{L}_n(\lambda) = \lim_{T \to \infty} \mathbb{E}\left[\sum_{t=0}^{T-1} \gamma^t r_n(s_n(t), a_n(t)) - \lambda a_n(t)\right]$$

- S(λ) : set of states to be idle obtained via L<sub>n</sub>(λ)
- Optimal policy : play (a = 1) if  $s \in S^{c}(\lambda)$  else idle (a = 0)
- Indexability : if idle for λ, then still idle for λ' > λ (if higher penalty for being active, stay idle)

#### Definition

$$\mathcal{IW}(S) = \lambda^*$$

s.t. 
$$S \in \partial \mathcal{P}(\lambda^*)$$
  
**Remark :** if  $\lambda < \lambda^*, S \in \mathcal{P}(\lambda)$ , else  
 $S \in \mathcal{P}^c(\lambda)$ 



Main result [unpublished]

We have  $S_u = X_u$ , and

$$\mathcal{IW}_u(X_u) = \sqrt{p_u}X_u$$

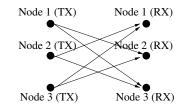
- Is it close to square-root law?
- Extension to 2-D state (age, popularity) when time-varying popularity
  - o to be complete
  - condition for indexability
  - never done for any 2-D state contex

## Perspective 1 : Wireless federative learning

- Learn a w-NN
- <u>but</u> database is split over K agents :  $(x_k, y_k)_{k=1, \dots, K}$  $w^* = \arg \min_{w} \sum_{k=1}^{K} f_k(w)$
- Agents are wirelessly connected Algorithm :
  - Local gradient computation :  $\nabla f_k(w_t)$
  - Sharing gradient and update :  $w_{t+1} \leftarrow w_t \mu \sum_{k=1}^{K} \nabla f_k(w_t)$

### Sharing step is a bottleneck!

- If no interference (baseline) : user rate =  $log_2(SNR)$
- If time-sharing : user rate =  $\frac{1}{K} \log_2(SNR)$
- If our scheme : user rate =  $\frac{K(K-1)-1}{K(2K-3)} \log_2(SNR) \sim \frac{1}{2} \log_2(SNR)$ More than **half the cake** for each agent ! (if K = 3, then 5/6)



## Perspective 1 : why does it work?

#### When you precode well, you have more than the cake :

• Example : interference channel

$$y_n = x_n + i \left(\sum_{m \neq n} x_m\right) + z_n, \ \forall n$$

• If 
$$x_n \in \mathbb{R}$$
, then  $\Re\{y_n\} = x_n + \Re\{z_n\}, \ \forall n$ 

User rate is then

$$\frac{1}{2}\log_2(1+\text{SNR})\sim \frac{1}{2}\log_2(\text{SNR})$$

 Interference Alignement (IA) : put interference in a common subspace at RX side. The orthogonal is interference-free

An other application : distributed estimation in complete graph (rather than star graph and 1-bit time-sharing communication) Collaboration with Shanghai Jiao Tong University (co-PhD degree) Representing an image (or set of patches) as a graph

#### Some examples :

- node = pixel; feature = color level; edge = two pixels belonging to similar patches; weight = location distance within the image
- node = patch; feature = patch learnt (strongly compressed) representation; edge+weight = related to original patch feature

Application : lung cancer

Collaboration with Universidad Nacional de Colombia (UNAL)

# Perspective 3 : Efficiency

 $\label{eq:Efficiency} \text{Efficiency} = \frac{\text{metric of performance}}{\text{consumed energy}}$ 

- OPEX-like energy : operational one
- CAPEX-like energy : embodied one
  - $\circ$  mining, manufacturing, recycling,  $\cdots$ : Life-Cycle Assessment

#### Example : Machine Learning

- Operational energy : computation energy during usage phase
- Embodied energy : Training, Computer's manufacturing, Cooling
- Metric of performance : Customer Satisfaction Rate

#### Main concerns :

- Open-data are missing for this kind of evaluation
- Depreciation duration, energy assignement (training/test)

Sustainable (or resilient) system "meets the needs of present generations without compromizing the ability of future generations to meet their own needs" [Brundtland1987]

Implementation : given an application/usage, level of power is fixed

Why is it different from energy efficient system?

- rebound effect has to be taken into account.
- if gain in energy consumption comes from enablement effect, customer behavior has to be predicted

#### Main concerns :

- Does not depend only on engineers' answers
- Required Science and Technology Studies (STS)

## Perspective 3 : sustainable systems

#### Cars' traffic management :

- Given an area, amount of energy is limited per prefixed duration
- Speed limited to satisfy the energy constraint
- Avoid Stop-and-Go policy (consuming the whole budget once)
- Long-term policy is required to be smoother
  - at which spatial scale : road, county (but long-haul traffic ?), country
  - o at which time scale : day, week, year
  - traffic prediction or adaptation?
- Machine Learning is a relevant tool since highly-complex problem

### Back to communication network :

- Given an area, amount of energy is limited per pre-fixed duration
- Packet traffic has to adapt
  - Quality of Service is moving
  - Outage is possible
- Here : available traffic model via stochastic geometry (ANR and PEPR grants in collaboration with *Infres*)