

# An optimal algorithm for resource allocation problem in concave context

A. Le Poupon and O. Rioul  
 ENST, 46, rue Barrault, 75634 Paris cedex 13  
 (Axel.LePoupon, Olivier.Rioul)@enst.fr

Keywords : Trees, Discrete Lagrangian Method, Dynamic programming

In this communication, an optimal algorithm is presented for solving the following resource allocation problem

$$\max \left\{ A = \sum_{i=1}^N A_i \mid B = \sum_{i=1}^N B_i \leq B_0 \right\} \quad (1)$$

where the  $(A_i, B_i)$  are constrained to a finite set of values and where  $B_0$  is the resource budget. We restrict our study to sets  $\{(A_i, B_i)\}$  such that  $A_i$  is a concave function of  $B_i$  (see [3] for a generalization). Applications include optimal bit allocation procedures for source coding [1] and optimal power allocation for multicarrier channel coding [2].

Up to now, (1) has not been solved completely by an efficient algorithm. Dynamic programming has disastrous computational complexity, while Lagrangian methods [1] give sub-optimal results because the optimal solution does not necessarily lie on the convex hull in the  $A$ - $B$  plane (see fig.1a).

Define a critical multiplier  $\lambda$  as one of the values  $\lambda = \Delta A / \Delta B$  for which two consecutive points on the convex hull are joined by a segment of slope  $\lambda$ . The Shoham-Gersho procedure [1] finds a point on the convex hull as the path corresponding to the whole sequence of multipliers  $\lambda_1 \geq \lambda_2 \cdots \geq \lambda_M$ . We propose a procedure to find the optimal solution to (1) obtained from a concave path corresponding to a subsequence  $\lambda_{i_1} \geq \lambda_{i_2} \cdots \geq \lambda_{i_m}$  (see fig.1b). The search of an inaccessible solution is made efficient by the means of a test criterion ensuring that all subsequent subpaths passing through a given point consist only of suboptimal points. They are therefore removed from the search which continues with surviving paths.

We obtain the following results : for  $N = 16$  sets of 16 values  $(A_i, B_i)$ , solving (1) with  $B_0$  covering the whole range (1000 test budgets), we find 100% of the optimal points when [1] finds in mean barely 5%, and we operate in 5.0 seconds when dynamic programming lasts 451.5 seconds (on Pentium4 @2.6GHz).

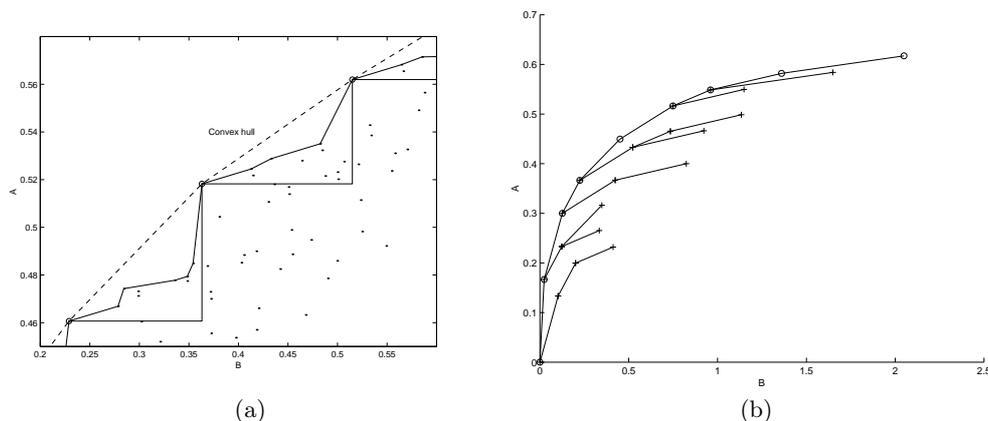


Figure 1: (a) Zoom on a  $A$ - $B$  set showing inaccessible solutions (inside triangles) between consecutive (circled) convex hull solutions. (b) Examples of concave paths drawn from origin  $(A, B) = (0, 0)$ .

[1] Y. Shoham, A. Gercho, *ASSP* 36, 1988, 1445-1453.

[2] B. S. Krongold *et al.*, *ITC* 48, 2000, 23-27.

[3] A. Le Poupon, O. Rioul, "On optimal resource allocation", *Trans. IT*, submitted, 2005