

Universal Coordinate Descent

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Introduction

Problem

Minimize a convex composite function

$$\min_{x \in \mathbb{R}^d} [F(x) = f(x) + \psi(x)]$$

Hölder gradient

We say that f has Hölder (sub)gradient ($f \in \text{Hölder}(M, \nu)$) if

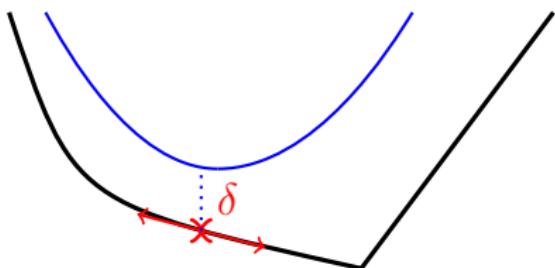
$$\|\nabla f(x) - \nabla f(y)\|_* \leq M \|x - y\|^\nu, \quad \forall x, y \in \mathbb{R}^d$$

- $\nu = 0$: non-differentiable but bounded subgradients
- $\nu = 1$: differentiable with Lipschitz gradient

Universal gradient methods

- Yu. Nesterov. Universal gradient methods for convex optimization problems, CORE DP #26/2013
- If $f \in \text{Hölder}(\mathcal{M}, \nu)$, then for all $\delta > 0$ and $L \geq \left(\frac{1}{\delta}\right)^{\frac{1-\nu}{1+\nu}} \mathcal{M}^{\frac{2}{1+\nu}}$:

$$f(x) \leq f(y) + \langle f(y), x - y \rangle + \frac{L}{2} \|x - y\|^2 + \frac{\delta}{2}$$



- ν “discovered” via a line search

Block coordinate setting

The Problem

$$\min_{x=(x^{(1)}, \dots, x^{(n)}) \in \mathbb{R}^d} [F(x) = f(x) + \psi(x)]$$

Blocks: $x^{(i)} \in \mathbb{R}^{n_i}$, $\sum_{i=1}^n n_i = d$

Assumptions

- **∇f is block-Hölder:** For each block $i \in \{1, 2, \dots, n\}$ and all $x \in \mathbb{R}^d$ and $h \in \mathbb{R}^{n_i}$,

$$\|\nabla_i f(x + U_i h) - \nabla_i f(x)\| \leq M_i \|h\|^{\nu}$$

- **ψ is block-separable:** $\psi(x) = \sum_{i=1}^n \psi_i(x^{(i)})$

Line search

Proposition

If $f \in \text{Hölder}(M, \nu)$, then for any $\delta > 0$ and $L \geq (\frac{1}{\delta})^{\frac{1-\nu}{1+\nu}} M^{\frac{2}{1+\nu}}$ and all $x, y \in \mathbb{R}^d$:

$$(1) \quad f(x) \leq f(y) + \langle f(y), x - y \rangle + \frac{L}{2} \|x - y\|^2 + \frac{\delta}{2}$$

$$(2) \quad f(y) + \langle \nabla f(y), x - y \rangle + \frac{1}{2L} \|\nabla f(x) - \nabla f(y)\|_*^2 - \frac{\delta}{2} \leq f(x)$$

$$(3) \quad \frac{1}{2L} \|\nabla f(x) - \nabla f(y)\|_*^2 \leq \frac{L}{2} \|x - y\|^2 + \delta$$

- In fact, $(1) \implies (2) \implies (3)$
- Moreover, if (3) holds for some L, x, y, δ , then

$$f(x) \leq f(y) + \langle \nabla f(y), x - y \rangle + L \|x - y\|^2 + \delta \quad (*)$$

Line search: based on “ $(3) \implies (*)$ ” for $f(h) \leftarrow f(x_k + U_i h)$

Universal Coordinate Descent (UCD)

Choose $x_0 \in \text{dom } \psi$, $L_0^1, \dots, L_0^n > 0$, $\epsilon > 0$ and $\delta = \frac{\epsilon}{2n}$

For $k \geq 0$ do:

1. **Select block:** $i \in \{1, 2, \dots, n\}$ at random
2. **Line search:** Find the **smallest integer s** such that for

$$h \leftarrow \arg \min_{h \in \mathbb{R}^{n_i}} f(x_k) + \langle \nabla_i f(x_k), h \rangle + 2^{s-1} L_k^i \|h\|^2 + \psi_i(x_k + h)$$

and $x_+ := x_k + U_i h$, we have

$$\frac{1}{2} \frac{1}{2^{s-1} L_k^i} \|\nabla_i f(x_k) - \nabla_i f(x_+)\|^2 \leq \frac{2^{s-1} L_k^i}{2} \|x_k^{(i)} - x_+^{(i)}\|^2 + \delta$$

3. **Update:** Set $x_{k+1} \leftarrow x_+$ and $L_{k+1}^i \leftarrow 2^s L_k^i$

Complexity

Theorem

The universal coordinate descent method satisfies:

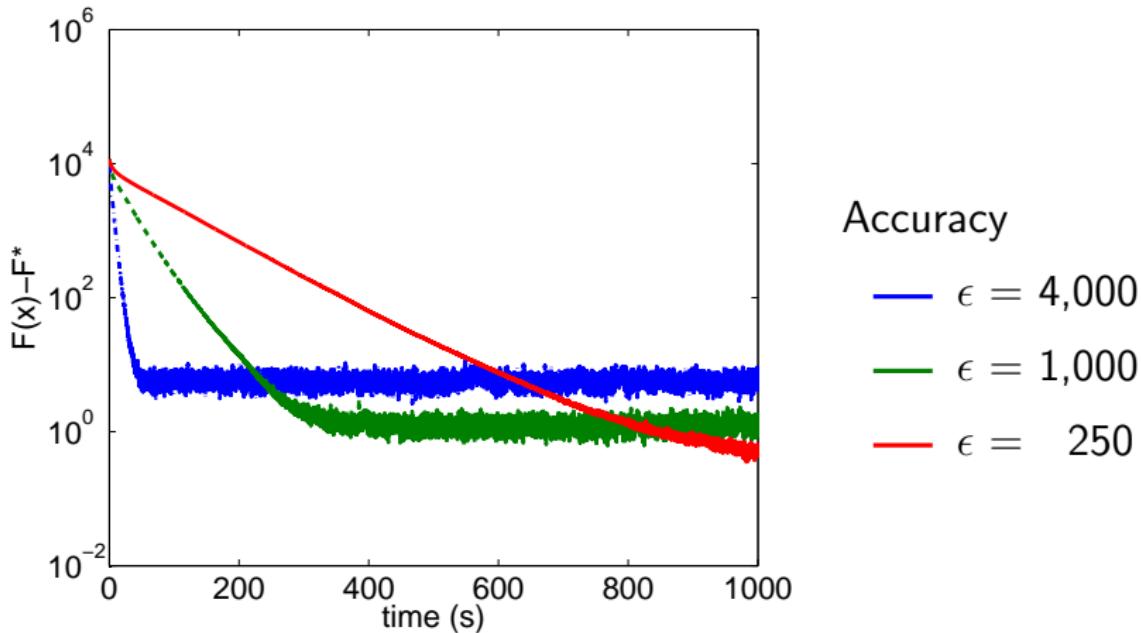
$$\begin{aligned} \mathbf{E}[F(\hat{x}_k) - F(x_*)] &\leq \frac{n}{k} R_0^2(x_*) + \frac{4n}{k} \bar{\Lambda} \mathcal{R}_M(x_*)^2 + \frac{\epsilon}{2} \\ &\leq \frac{n}{k} R_0^2(x_*) + \frac{4n^{\frac{2}{1+\nu}}}{k \epsilon^{\frac{1-\nu}{1+\nu}}} \bar{\Lambda} \mathcal{R}_{M^{\frac{2}{1+\nu}}}^2(x_*) + \frac{\epsilon}{2}, \end{aligned}$$

where $\mathcal{R}_M(x_*) = \sup_k \|x_k - x_*\|_M$ and $\bar{\Lambda} = \frac{1}{n} \sum_{i=1}^n \log_2\left(\frac{2M^j}{L_0^i}\right)$.
 \Rightarrow # iterations needed to reach an ϵ -solution

- $\nu = 0$: $O(n^2/\epsilon^2)$
- $\nu = 1$: $O(n/\epsilon)$

Choice of ϵ

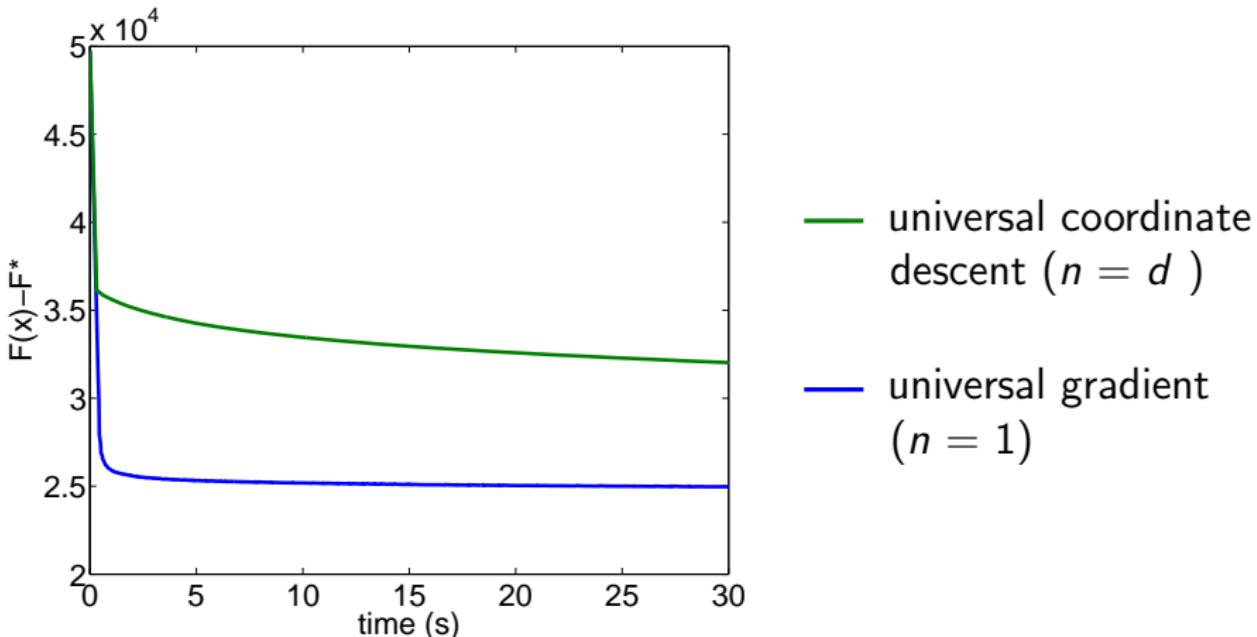
w8a dataset: $d = 300$, $m = 49749$. $F(x) = \|Ax - b\|_1 + \|x\|_1$



No convergence but ϵ -solutions

Subgradient-based coordinate descent

w8a dataset: $d = 300$, $m = 49749$. $F(x) = \|Ax - b\|_1 + \|x\|_1$



Coordinate descent not so efficient in this case: $O(n^2/\epsilon^2)$

Parallel coordinate descent

Studied in:

- P. Richtárik and M. Takáč. Parallel coordinate descent methods for big data optimization. *Mathematical Programming* 144(2):1–38, 2014.
- O. Fercoq and P. Richtárik. Smooth minimization of nonsmooth functions with parallel coordinate descent methods. *arXiv:1309.5885*, 2013
- P. Richtárik and M. Takáč. Distributed coordinate descent method for learning with big data. *arXiv:1310.2059*, 2013
- O. Fercoq and P. Richtárik. Accelerated, parallel and proximal coordinate descent. *arXiv:1312.5799*, 2013

Parallel coordinate descent

- Update several coordinates at the same time:
random sampling \hat{S}
- **Theory** based on the concept of Expected Separable Overapproximation (ESO):

$$\mathbf{E} \left[\mathbf{f}(x + h_{[\hat{S}]}) \right] \leq \mathbf{f}(x) + \frac{\mathbf{E}[|\hat{S}|]}{n} \left(\langle \nabla \mathbf{f}(x), h \rangle + \frac{1}{2} \|h\|_v^2 \right)$$

- Vector $v = (v_1, \dots, v_n)$ depends on \hat{S} and f
- The dependence of v on

$$\tau := \mathbf{E}[|\hat{S}|]$$

determines the **parallelization speedup**

Partial separability

Definition

Assume $f : \mathbb{R}^d \rightarrow \mathbb{R}$ is **partially separable**. That is,

$$f(x) = \sum_{j=1}^m f_j(x),$$

where f_j depends on blocks $i \in C_j \subseteq \{1, 2, \dots, n\}$ only and

$$\omega = \max_j |C_j| \quad (\text{degree of partial separability})$$

- Note that $1 \leq \omega \leq n$

Deterministic Separable Overapproximation

Proposition (Non-quadratic DSO)

If $f : \mathbb{R}^d \rightarrow \mathbb{R}$ is convex and partially separable of degree ω , then for any $S \subseteq \{1, \dots, d\}$ and $x, h \in \mathbb{R}^d$,

$$f \left(x + \sum_{i \in S} U_i h^{(i)} \right) - f(x) \leq \frac{1}{\omega_S} \sum_{i \in S} (f(x + \omega_S U_i h^{(i)}) - f(x)),$$

where $\omega_S = \max_j |C_j \cap S| \leq \min(\omega, |S|)$

Line search:

Estimate the curvature of $\phi_i(h^{(i)}) = f(x + \omega_S U_i h^{(i)}) - f(x)$ independently for all $i \in S$

Expected Separable Overapproximation

Proposition (Non-quadratic ESO)

Let $f : \mathbb{R}^d \rightarrow \mathbb{R}$ be convex and partially separable of degree ω and \hat{S} be a τ -nice sampling and let

$$\pi_j = \binom{\omega}{j} \binom{n - \omega}{\tau - j} / \binom{n}{\tau}$$

For all $x, h \in \mathbb{R}^d$:

$$\mathbf{E}[f(x + h_{[\hat{S}]}) - f(x)] \leq \sum_{i=1}^n \sum_{j=1}^{\min(\omega, \tau)} \frac{\pi_j}{\omega} (f(x + jU_i h^{(i)}) - f(x))$$

Simpler ESO

Proposition

*The ESO can be further bounded by the following separable function involving a **logarithmic number of terms per block**:*

$$\mathbf{E}[f(x + h_{[\hat{S}]}) - f(x)] \leq \frac{\tau}{n} \sum_{i=1}^n \tilde{f}_x^i(h^{(i)}),$$

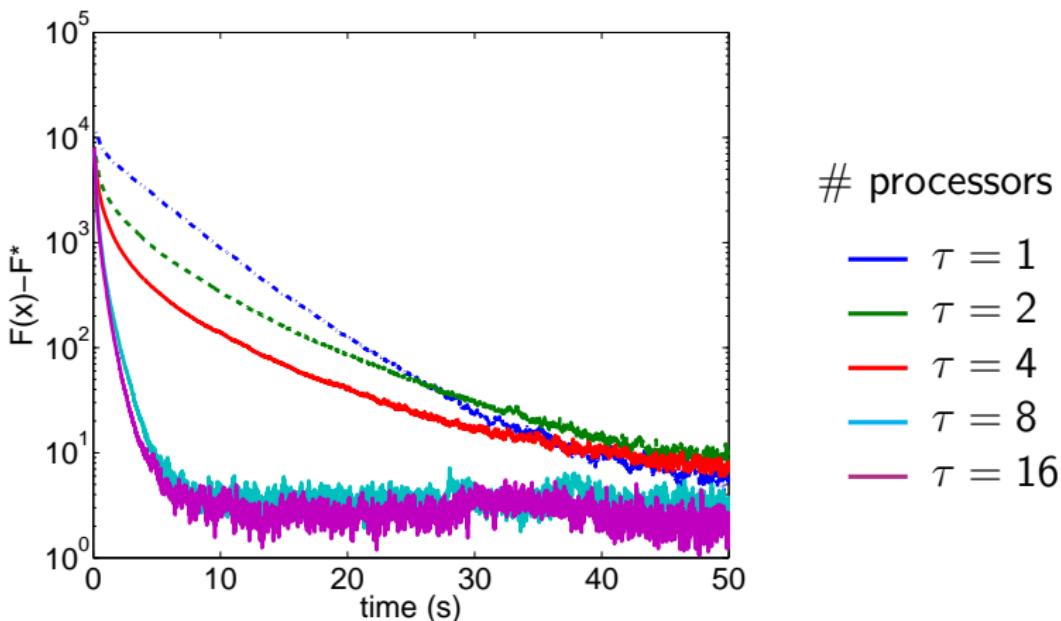
where

$$\begin{aligned} \tilde{f}_x^i(h) &:= \sum_{s=0}^{\lfloor \log_2(\min(\omega, \tau)) \rfloor} \frac{Q_s}{2^{s+1}-1} (f(x + (2^{s+1}-1)U_i h^{(i)}) - f(x)) \\ Q_s &= \frac{n}{\omega \tau} \sum_{j=2^s}^{2^{s+1}-1} j \pi_j. \end{aligned}$$

Parallelization speedup

w8a dataset: $d = n = 300$, $m = 49,749$

$$F(x) = \|Ax - b\|_1 + \|x\|_1, \quad \epsilon = 4,000$$



Acceleration

- For smooth functions rate $O(1/\sqrt{\epsilon})$ instead of $O(1/\epsilon)$
- If nonsmooth: no improvement, rate remains $O(1/\epsilon^2)$
- Arguments of line search and acceleration combine well
- Algorithm a little more complex

Universal Accelerated Coordinate Descent

Choose $x_0 \in \text{dom}\psi$; $L_0^1, \dots, L_0^n > 0$; $\epsilon > 0$. Set $z_0 = x_0$ and $\theta_0 = \frac{\tau}{n}$.

For $k \geq 0$ do:

1. $y_k = (1 - \theta_k)x_k + \theta_k z_k$
2. Pick a subset of τ blocks \hat{S} , uniformly at random
3. In parallel for $i \in \hat{S}$:
 - 3.1 Find the smallest integer s such that for

$$z_+^{(i)} = \arg \min_{z \in \mathbb{R}^{n_i}} \langle \nabla_i f(y_k), z \rangle + \frac{n\theta_k 2^s L_k^i}{2\tau} \|z - z_k^{(i)}\|^2 + \psi_i(z)$$

$$h_k^{(i)} = \frac{n\theta_k}{\tau} (z_+^{(i)} - z_k^{(i)})$$

we have $\frac{1}{2} \frac{1}{2^{s-1} L_k^i} \|\nabla_i f(y_k) - \nabla_i \tilde{f}_{y_k}^{(i)}(h_k^{(i)})\|^2 \leq \frac{2^{s-1} L_k^i}{2} \|h_k^{(i)}\|^2 + \frac{\epsilon \theta_k}{2\tau}$

3.2 Set $L_{k+1}^i = 2^s L_k^i$ and $x_{k+1}^{(i)} = y_k^{(i)} + h_k^{(i)}$

$$4. \theta_{k+1} = \frac{\sqrt{\theta_k^4 + 4\theta_k^2} - \theta_k^2}{2}$$

Complexity

Theorem

The universal accelerated coordinate descent method satisfies

$$\begin{aligned} \mathbf{E}[F(x_{k+1}) - F(x_*)] &\leq \frac{4n^2}{(\textcolor{blue}{k}\tau + 2n)^2} \left(\left(1 - \frac{\tau}{n}\right) (F(x_0) - F(x_*)) + \frac{1}{2} R_{L_0}^2(x_*) \right) \\ &\quad + \left(\frac{2n}{\textcolor{blue}{k}\tau + 2n} \right)^{\frac{1+3\nu}{1+\nu}} \frac{(2n)^{\frac{1-\nu}{1+\nu}} \beta_\nu C}{\epsilon^{\frac{1-\nu}{1+\nu}}} \mathcal{R}_{M^{\frac{2}{1+\nu}}}^2(x_*) + \frac{\epsilon}{2} \end{aligned}$$

where $\mathcal{R}_M(x_*) = \sup_k \|x_k - x_*\|_M$, $C = 2 \left(\bar{\Lambda} + \frac{1-\nu}{1+\nu} \log\left(\frac{\textcolor{blue}{k}\tau + 2n}{\epsilon}\right) \right)$

$$\beta_\nu = \sum_{j=1}^n \frac{n\pi_j}{\tau\omega} j^{1+\nu}, \quad \bar{\Lambda} = \frac{\tau}{n} \sum_{i=1}^n \log_2 \frac{4\beta_\nu(M_i)^{\frac{2}{1+\nu}}}{L_0^i}$$

Accelerated Adaboost

- Semi-supervised classification
- A : matrix of characteristics
 b : labels
- Adaboost = greedy coordinate descent for

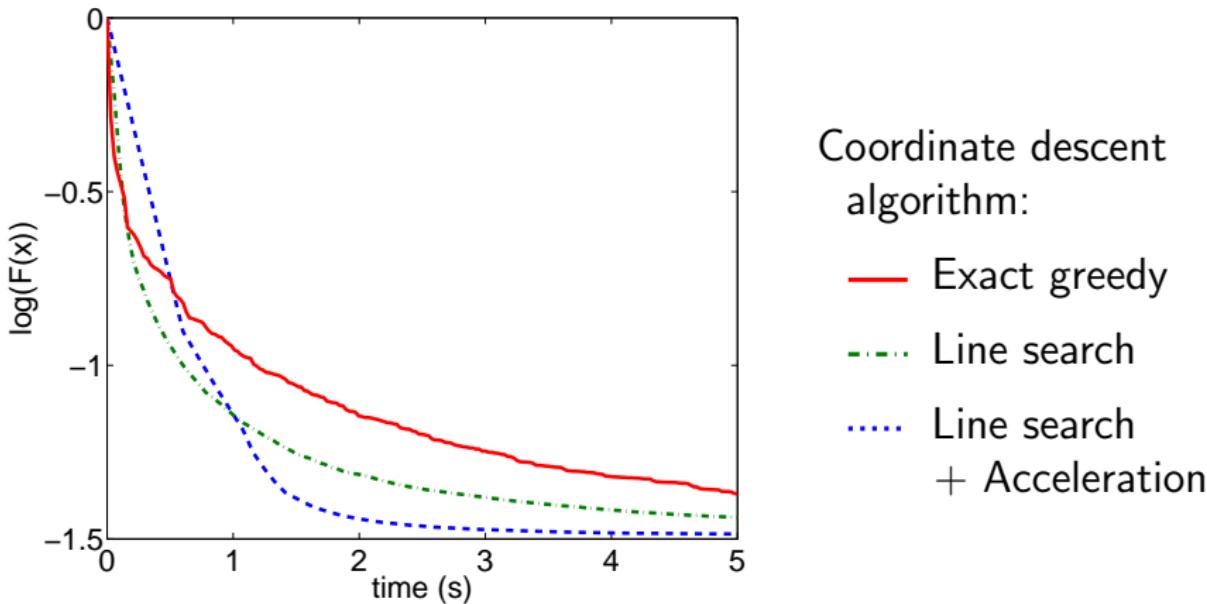
$$F(x) = \frac{1}{m} \sum_{j=1}^m \exp(b_j \sum_{i=1}^n A_{j,i} x_i)$$

- Differentiable but not globally Lipschitz

Numerical experiment

w8a dataset: $d = 300, m = 49749$

$$F(x) = \frac{1}{m} \sum_{j=1}^m \exp(b_j \sum_{i=1}^n A_{j,i} x_i), \epsilon = 0, 16 \text{ processors}$$



Conclusion

Summary:

- Universal coordinate descent algorithms:
accelerated, parallel, proximal
- Versatile line search
- Coordinates tackled independently

Open questions:

- Improve the theoretical bound on nonsmooth problems from $O(n^2/\epsilon^2)$ to $O(n/\epsilon^2)$
- Decrease the curvature estimate on the run
- Use non-quadratic ESO with exact minimization
(for smooth functions)